

Conservative Cascades: an Invariant of Internet Traffic

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Schedule

1) **Scaling theory**



Scaling
theory

2) **Cascades and TCP traffic**



Scaling
practice

3) **TCP traffic analysis**



Traffic
analysis

4) **Conclusions**

Scaling

~ absence of characteristic (time)scale

~ dependence among (time)scales

Many different types of scaling flavours :

- self-similarity
- long-range dependence
- multiscaling
- multifractality
- infinitely divisible cascades

Terminology

Let $X(t)$ be a regularly sampled process :

- $d(j,k)$: process increment at timescale j and time k
- q : moment of the process
- j : timescale, $j = 1, \dots, n$.
- $E|X(t)|^q$: expectation of moment q of process X at time t (over all possible realizations of the process).

Assumption :

- $d(j,k)$ is a homogeneous (stationary) zero-mean process.

Self-similarity

$$E|d(j,.)|^q = C_q (2^j)^{qH} \propto \exp(q H \ln(2^j))$$

- linear scaling among timescales and linear scaling among moments both driven by a single parameter H

⇒ called **mono-scaling** since H is constant

Multiscaling

$$E|d(j,.)|^q = C_q (2^j)^{H(q)} \propto \exp(H(q) \ln(2^j))$$

- linear scaling among timescales j , but driven by a function $H(\cdot)$ of the moment q
- \Rightarrow **mono-scaling becomes multi-scaling** since scaling depends on the moment q

Infinitely divisible cascade

$$E|d(j,.)|^q = C_q (2^j)^{H(q)} \propto \exp(H(q) n(2^j))$$

where

- $H(q)$ represents 1 step of the cascade
- $n(2^j)$ represents cascade depth at timescale j

⇒ scaling does not depend linearly on the moment q nor the timescale j anymore

The missing link

$$E|d(j, \cdot)|^q \propto \exp(f(q) g(2^j))$$

- **separability** of the moments (q) and the timescales (j)
 - higher-order moments emphasize on larger irregularities
 - larger timescales focus on smoothed versions of the process (zooming out)
- \Rightarrow framework to study irregularities and timescales independently**

Probabilistic interlude

- Cascade : multiplicative process that breaks a process into smaller and smaller fragments according to some (deterministic or random) rule.
- Link with infinitely divisible distributions (Feller vol. 2) :
F is infinitely divisible if for every n there exists a distribution F_n such that $F = F_n^{n*}$ where “*” denotes the convolution operator.

or equivalently

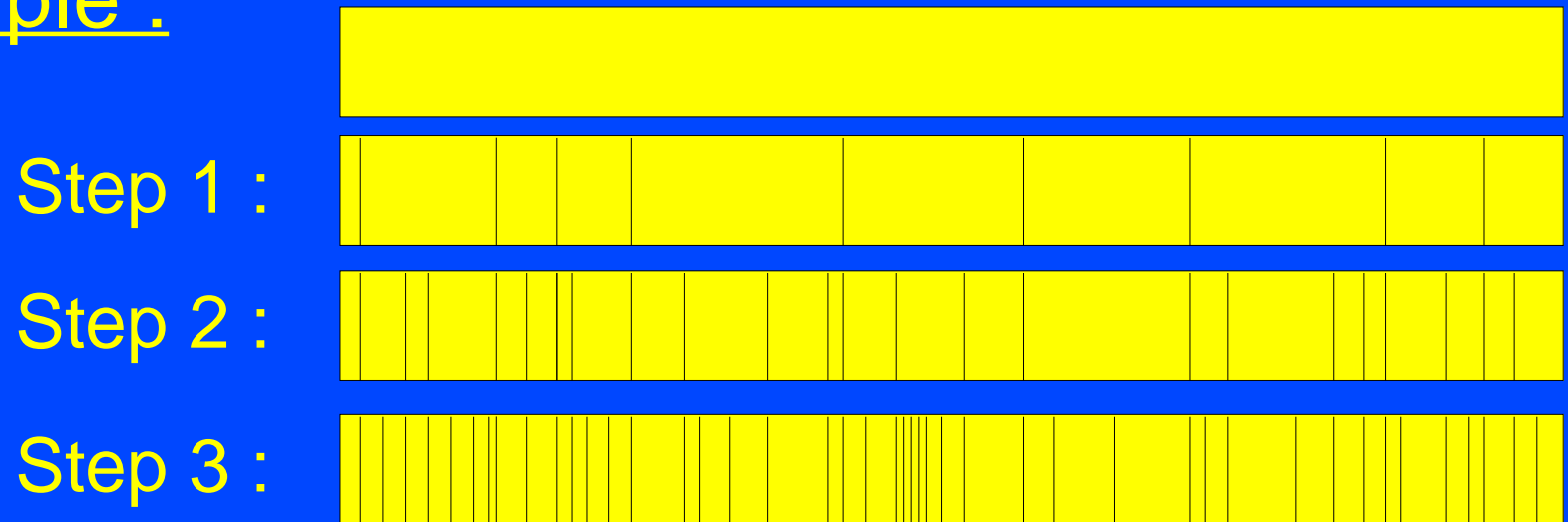
F is infinitely divisible iff for each n it can be represented as the distribution of the sum $S_n = X_{1,n} + \dots + X_{n,n}$ of n independent random variables with a common distribution F_n .

Conservative cascades

2 objectives :

- preservation of total mass of the process
- randomness in the splitting of the mass of the process

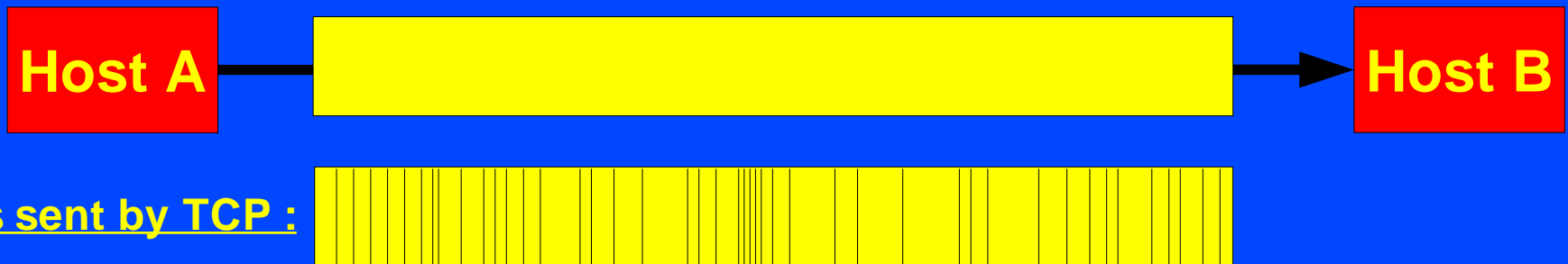
Example :



Conservative cascades and TCP traffic

TCP :

- breaks the total mass of the data to be sent into segments
- application sending data, TCP state machine and random network conditions drive how data segments are broken and when they are sent over time



Cascade parameters and TCP

- $H(q)$: - describes cascade generator (1 step of the breaking of the data)
- represents how TCP distributes the mass of the TCP segments into large and small irregularities
- $n(2^j)$: how many times the generator has been applied at timescale j (how many times the convolution operator has been applied)
- ⇒ $H()$ and $n()$ “*summarize*” the numerical behavior of the cascade

Wavelet analysis

- process increments $d(j,k)$ replaced by wavelet coefficients $\delta(j,k)$
- Wavelet-based partition function :

$$S(q,j) = \sum_k |2^{-j/2} \delta(j,k)|^q \quad (\propto E|d(j,k)|^q)$$

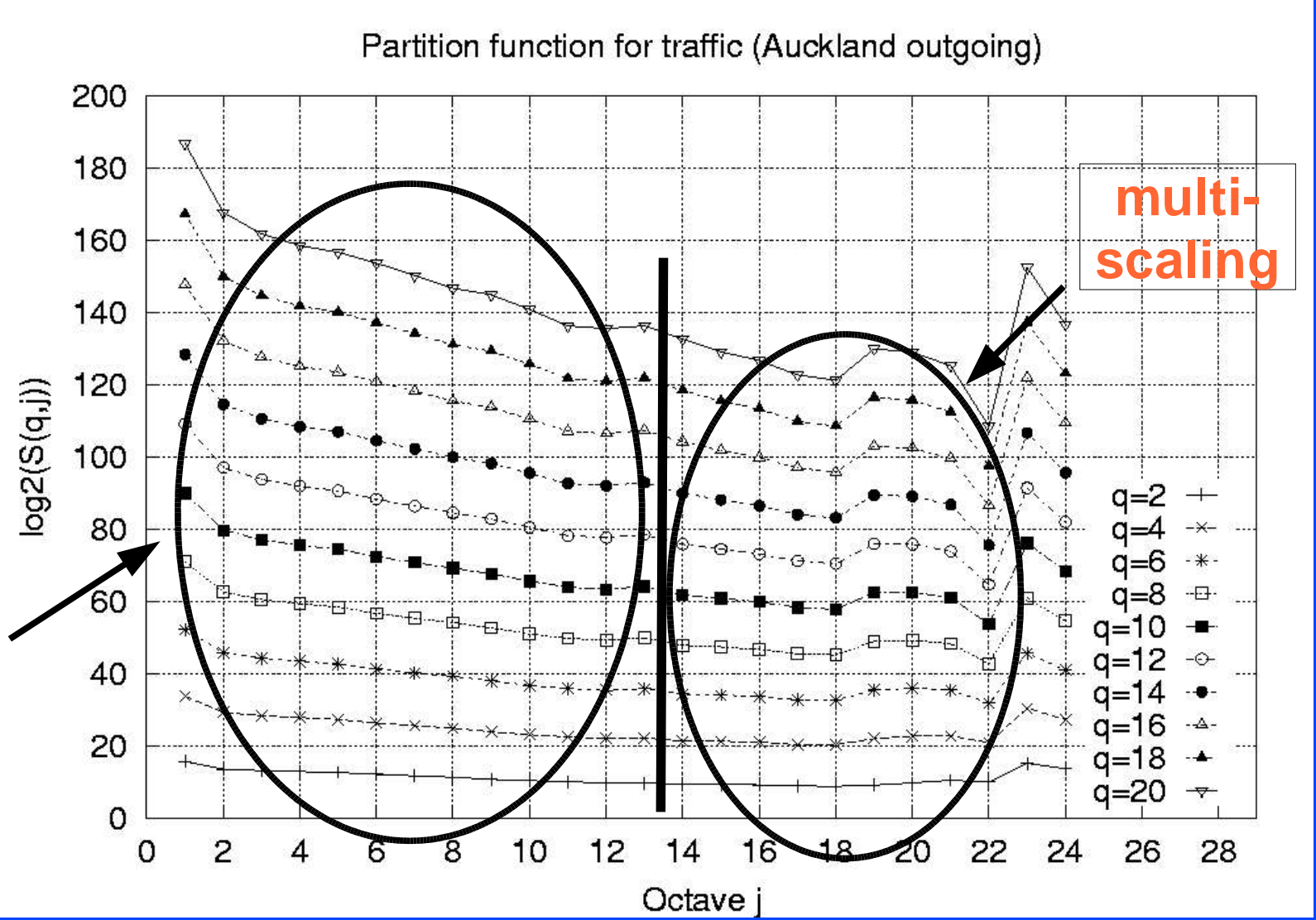
- Estimation of $H(q)$ and $n(2^j)$ via regression :

$$\ln(S(q,j)) \cong - H(q) \ln(2^j) + K_q$$

Internet traffic analysis

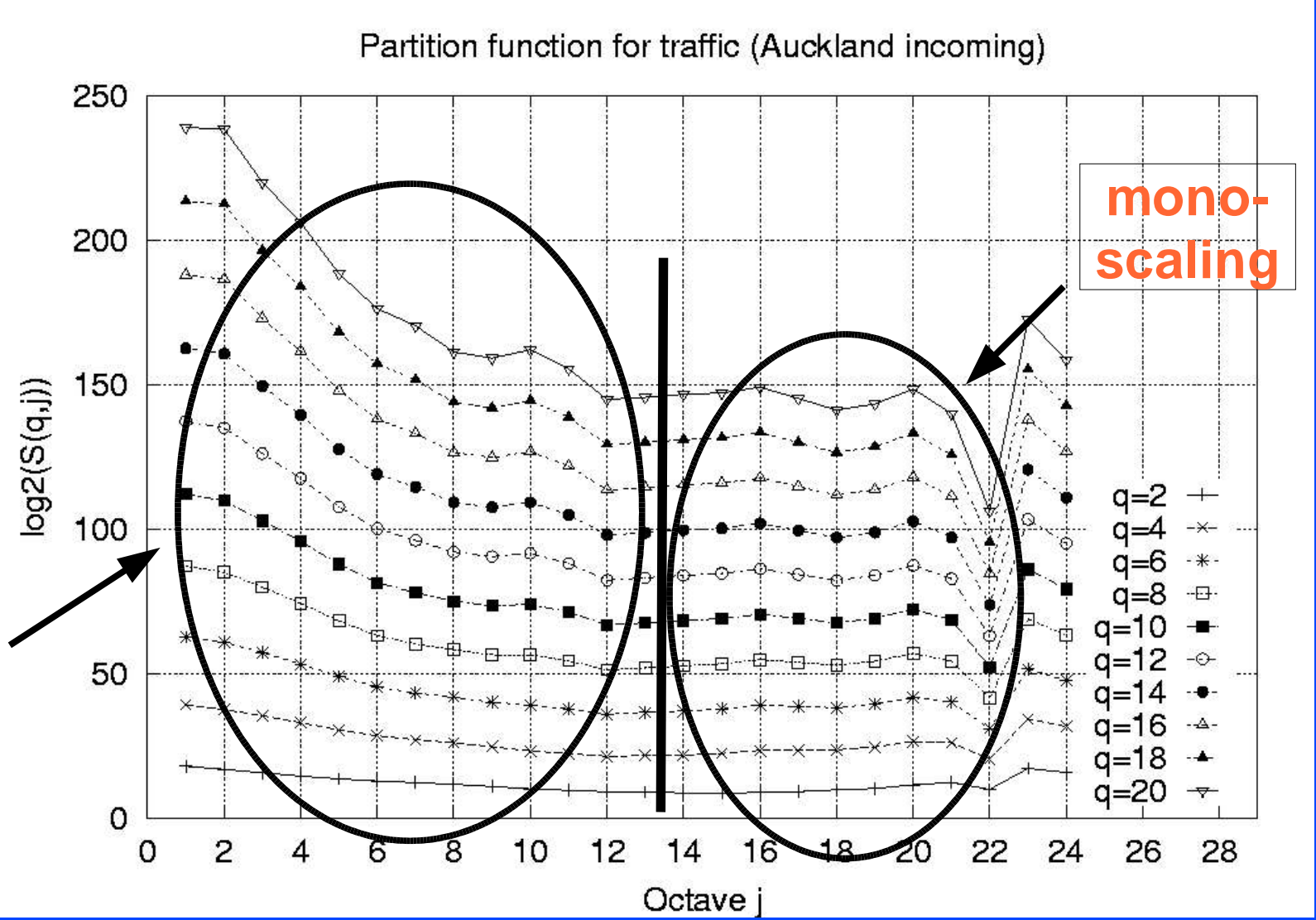
- 15 hours of IP packets from Auckland
- both traffic sent to the Internet and received from the Internet
- 1,629,069 incoming TCP flows, 1,613,976 outgoing
- more than 13 GBytes of incoming traffic, 10 GBytes outgoing
- 1 μ s precision, 1 ms time granularity used.

Multiscaling analysis (out)



multi-scaling

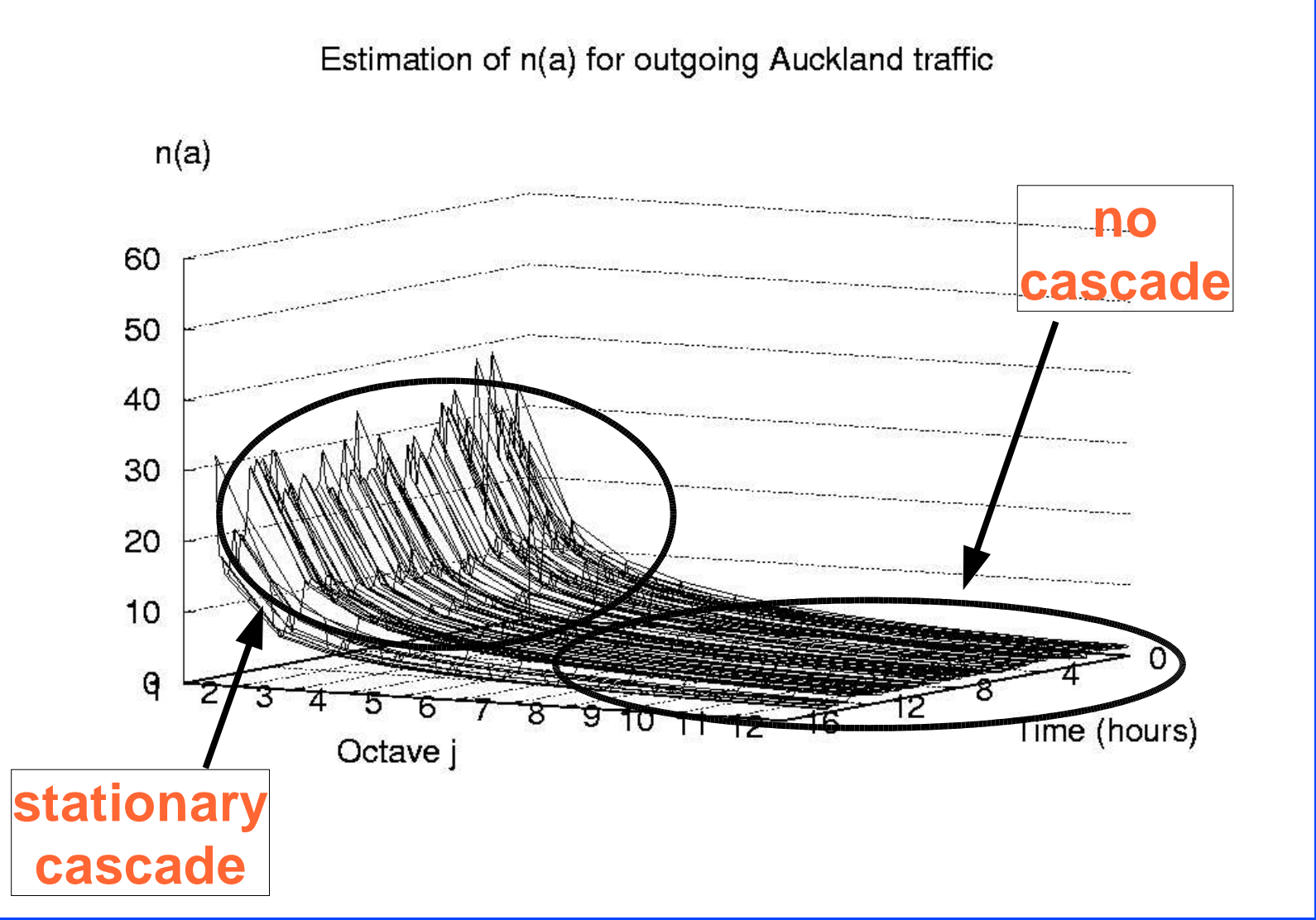
Multiscaling analysis (in)



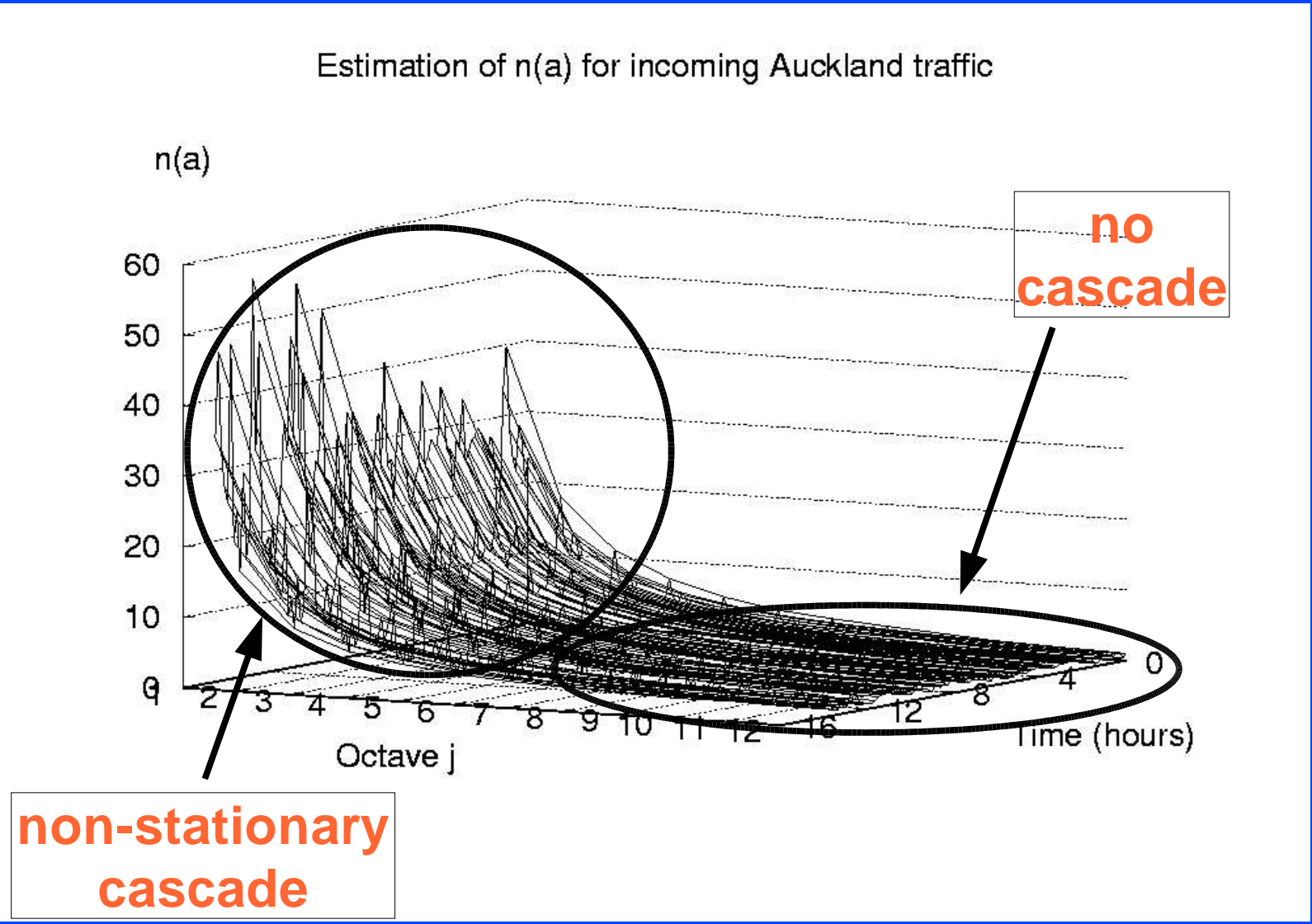
multi-scaling

mono-scaling

Cascade analysis (out)

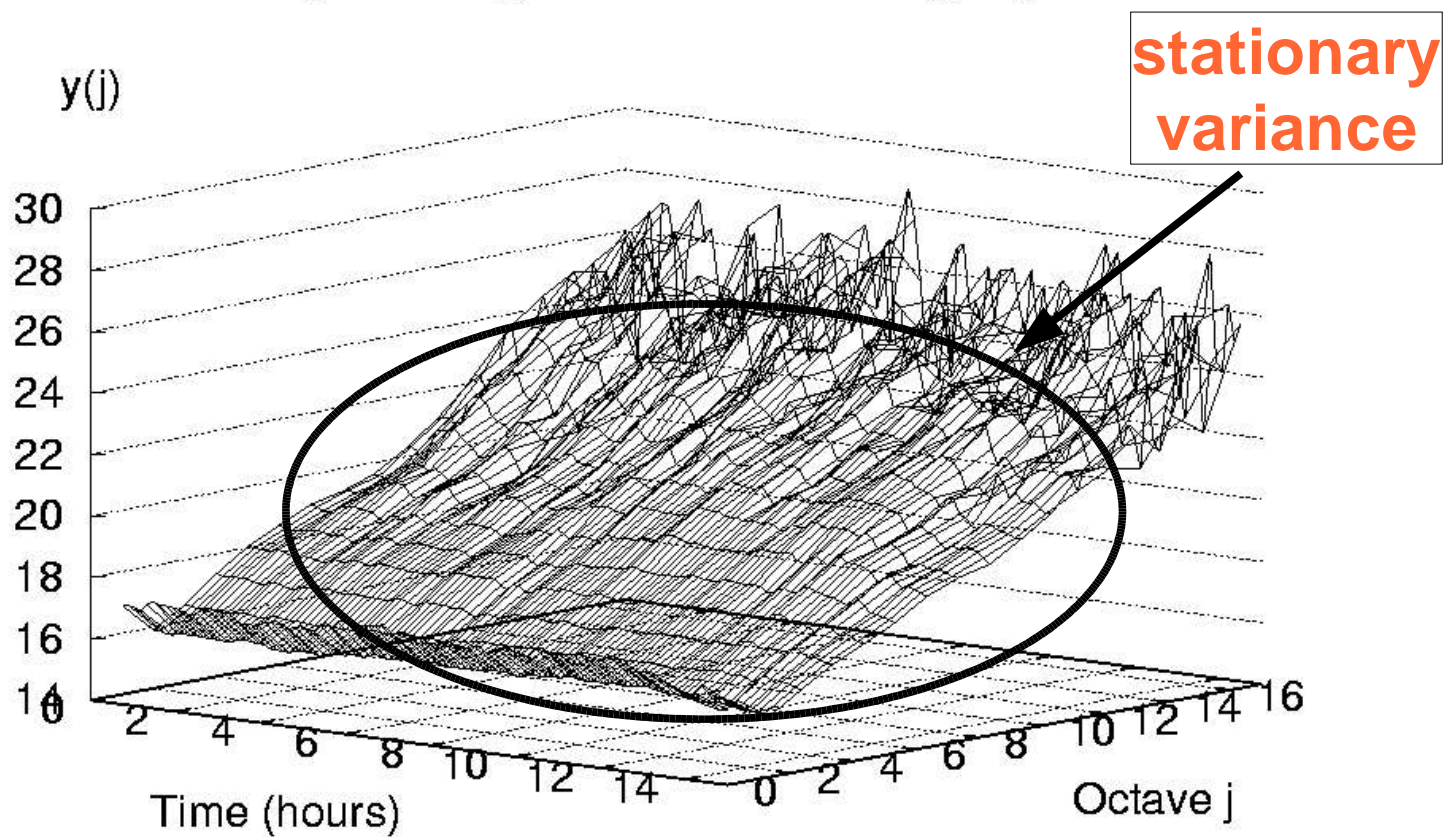


Cascade analysis (in)



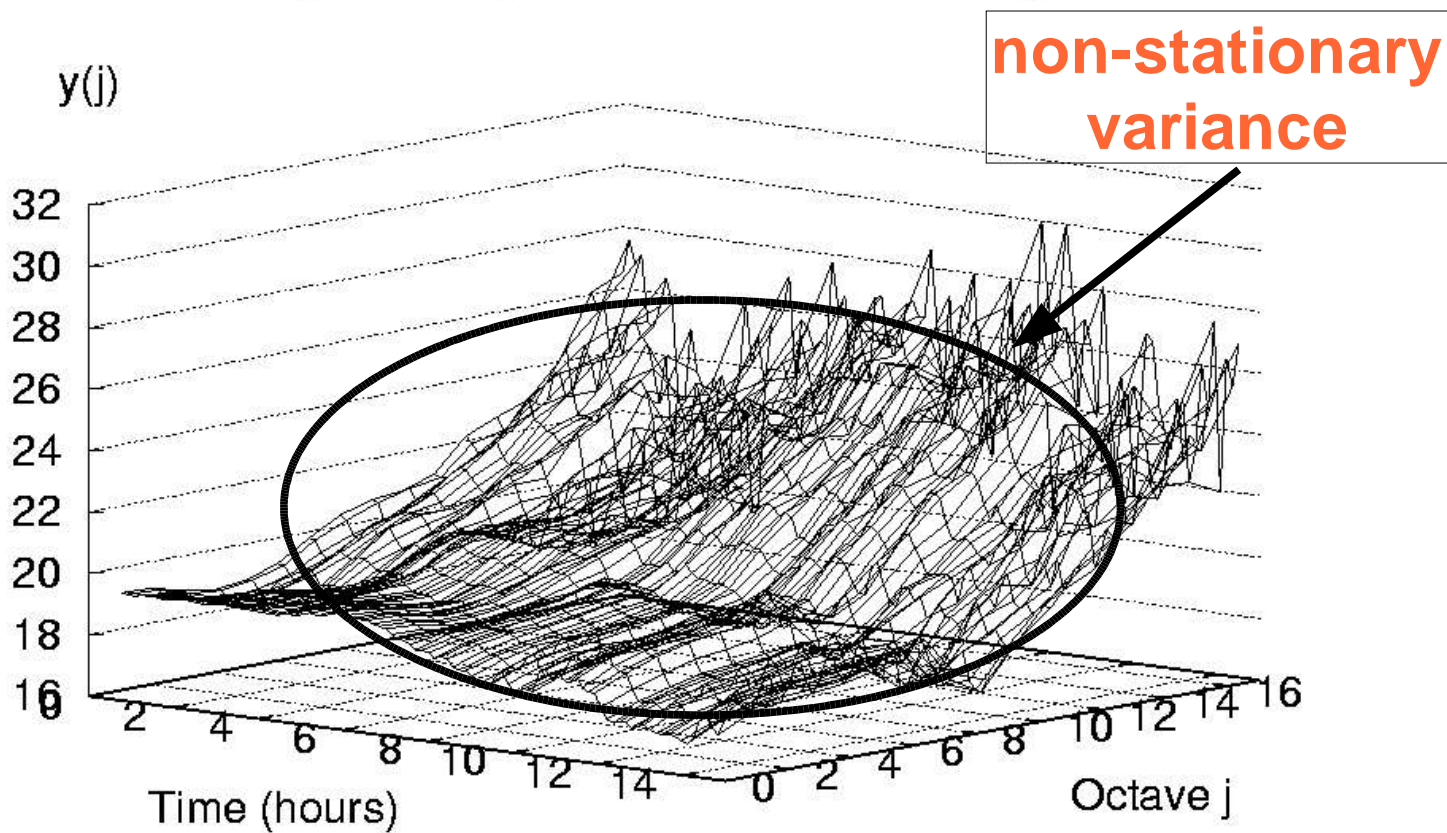
Variance analysis (out)

3D-logscale diagram for Auckland outgoing traffic



Variance analysis (in)

3D-logscale diagram for Auckland incoming traffic



Conclusions

- Cascade model captures well invariance in TCP behavior
- Cascade model does not fully capture traffic dynamics, only TCP traffic segmentation
- Cascade parameters do not show effect of cross-traffic, 2nd order properties do

After-thoughts

- Cascade is OK to generate clean TCP traffic at timescales below seconds
- Cascade is not OK to simulate realistic Internet traffic
- More work necessary to understand role of the “abstract” cascade parameters on the traffic behavior
- Can we distinguish between “normal” and “abnormal” TCP flows thanks to real-time monitoring of cascade parameters ?