



# Conservative Cascades: an Invariant of Internet Traffic

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# Schedule

**1) Scaling theory**



Scaling  
theory

**2) Cascades and TCP traffic**



Scaling  
practice

**3) TCP traffic analysis**



Traffic  
analysis

**4) Conclusions**

# Scaling

~ absence of characteristic (time)scale

~ dependence among (time)scales

Many different types of scaling flavours :

- self-similarity
- long-range dependence
- multiscaling
- multifractality
- infinitely divisible cascades

# Terminology

Let  $X(t)$  be a regularly sampled process :

- $d(j,k)$  : process increment at timescale  $j$  and time  $k$
- $q$  : moment of the process
- $j$  : timescale,  $j = 1, \dots, n$ .
- $E|X(t)|^q$  : expectation of moment  $q$  of process  $X$  at time  $t$  (over all possible realizations of the process).

Assumption :

- $d(j,k)$  is a homogeneous (stationary) zero-mean process.

# Self-similarity

$$E|d(j,.)|^q = C_q (2^j)^{qH} \propto \exp(q H \ln(2^j))$$

- linear scaling among timescales and linear scaling among moments both driven by a single parameter  $H$

⇒ called **mono-scaling** since  $H$  is constant

# Multiscaling

$$E|d(j,.)|^q = C_q (2^j)^{H(q)} \propto \exp(H(q) \ln(2^j))$$

- linear scaling among timescales  $j$ , but driven by a function  $H(\cdot)$  of the moment  $q$
- $\Rightarrow$  **mono-scaling becomes multi-scaling** since scaling depends on the moment  $q$

# Infinitely divisible cascade

$$E|d(j,.)|^q = C_q (2^j)^{H(q)} \propto \exp(H(q) n(2^j))$$

where

- $H(q)$  represents 1 step of the cascade
- $n(2^j)$  represents cascade depth at timescale  $j$

⇒ scaling does not depend linearly on the moment  $q$  nor the timescale  $j$  anymore

# The missing link

$$E|d(j, \cdot)|^q \propto \exp(f(q) g(2^j))$$

- **separability** of the moments ( $q$ ) and the timescales ( $j$ )
  - higher-order moments emphasize on larger irregularities
  - larger timescales focus on smoothed versions of the process (zooming out)
- $\Rightarrow$  framework to study irregularities and timescales independently**

# Probabilistic interlude

- Cascade : multiplicative process that breaks a process into smaller and smaller fragments according to some (deterministic or random) rule.
- Link with infinitely divisible distributions (Feller vol. 2) :  
F is infinitely divisible if for every n there exists a distribution  $F_n$  such that  $F = F_n^{n*}$  where “\*” denotes the convolution operator.

or equivalently

F is infinitely divisible iff for each n it can be represented as the distribution of the sum  $S_n = X_{1,n} + \dots + X_{n,n}$  of n independent random variables with a common distribution  $F_n$ .

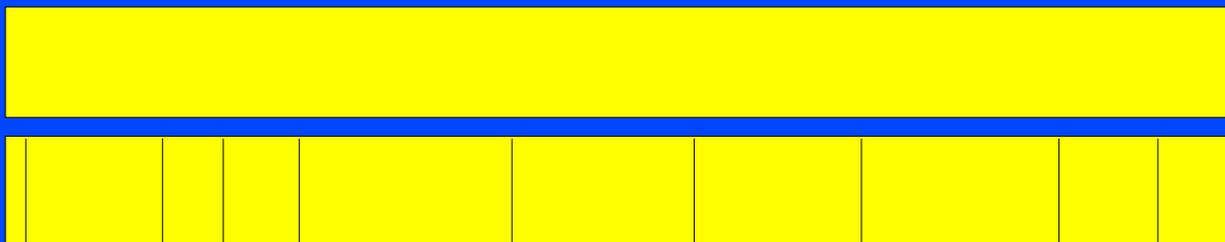
# Conservative cascades

## 2 objectives :

- preservation of total mass of the process
- randomness in the splitting of the mass of the process

## Example :

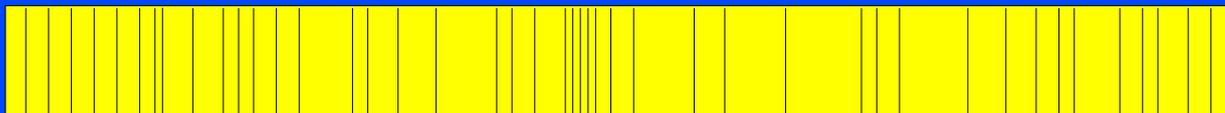
Step 1 :



Step 2 :



Step 3 :



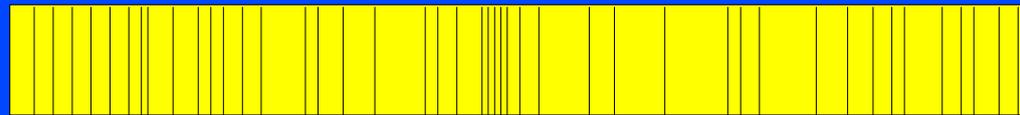
# Conservative cascades and TCP traffic

## TCP :

- breaks the total mass of the data to be sent into segments
- application sending data, TCP state machine and random network conditions drive how data segments are broken and when they are sent over time



Segments sent by TCP :



# Cascade parameters and TCP

- $H(q)$  : - describes cascade generator (1 step of the breaking of the data)
- represents how TCP distributes the mass of the TCP segments into large and small irregularities
- $n(2^j)$  : how many times the generator has been applied at timescale  $j$  (how many times the convolution operator has been applied)
- ⇒  $H()$  and  $n()$  “*summarize*” the numerical behavior of the cascade

# Wavelet analysis

- process increments  $d(j,k)$  replaced by wavelet coefficients  $\delta(j,k)$
- Wavelet-based partition function :

$$S(q,j) = \sum_k |2^{-j/2} \delta(j,k)|^q \quad (\propto E|d(j,k)|^q)$$

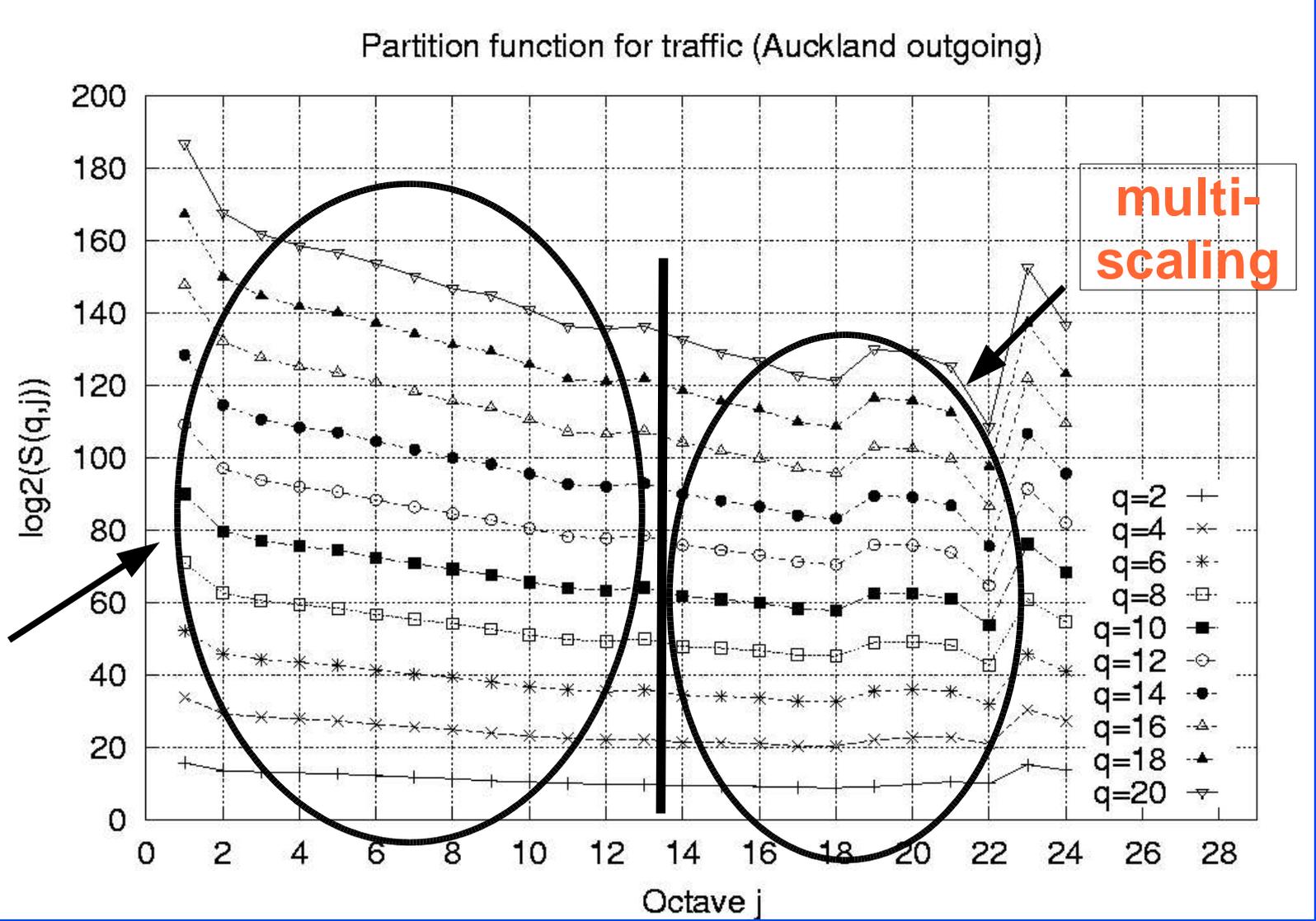
- Estimation of  $H(q)$  and  $n(2^j)$  via regression :

$$\ln(S(q,j)) \cong - H(q) \ln(2^j) + K_q$$

# Internet traffic analysis

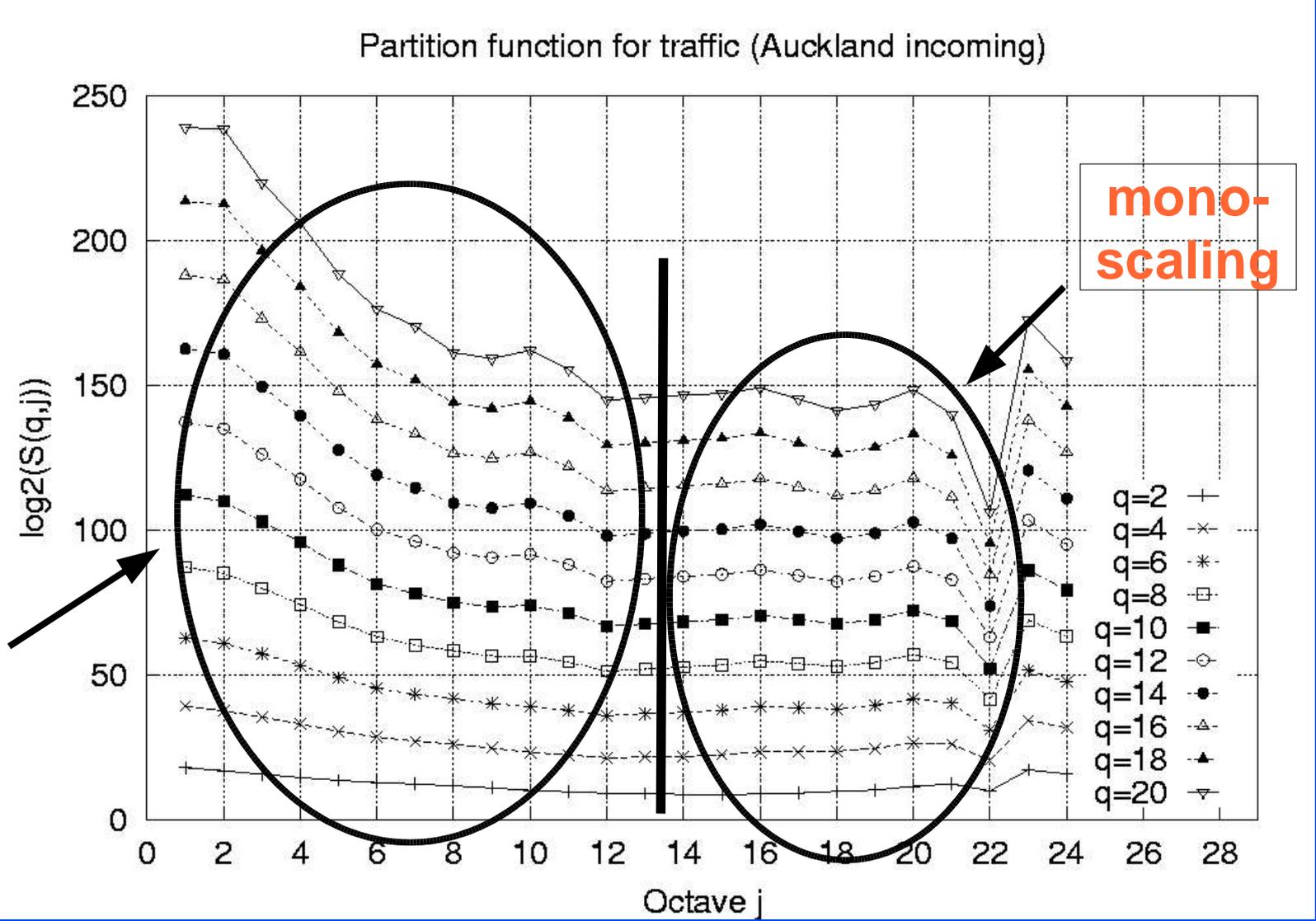
- 15 hours of IP packets from Auckland
- both traffic sent to the Internet and received from the Internet
- 1,629,069 incoming TCP flows, 1,613,976 outgoing
- more than 13 GBytes of incoming traffic, 10 GBytes outgoing
- 1  $\mu$ s precision, 1 ms time granularity used.

# Multiscaling analysis (out)



multi-scaling

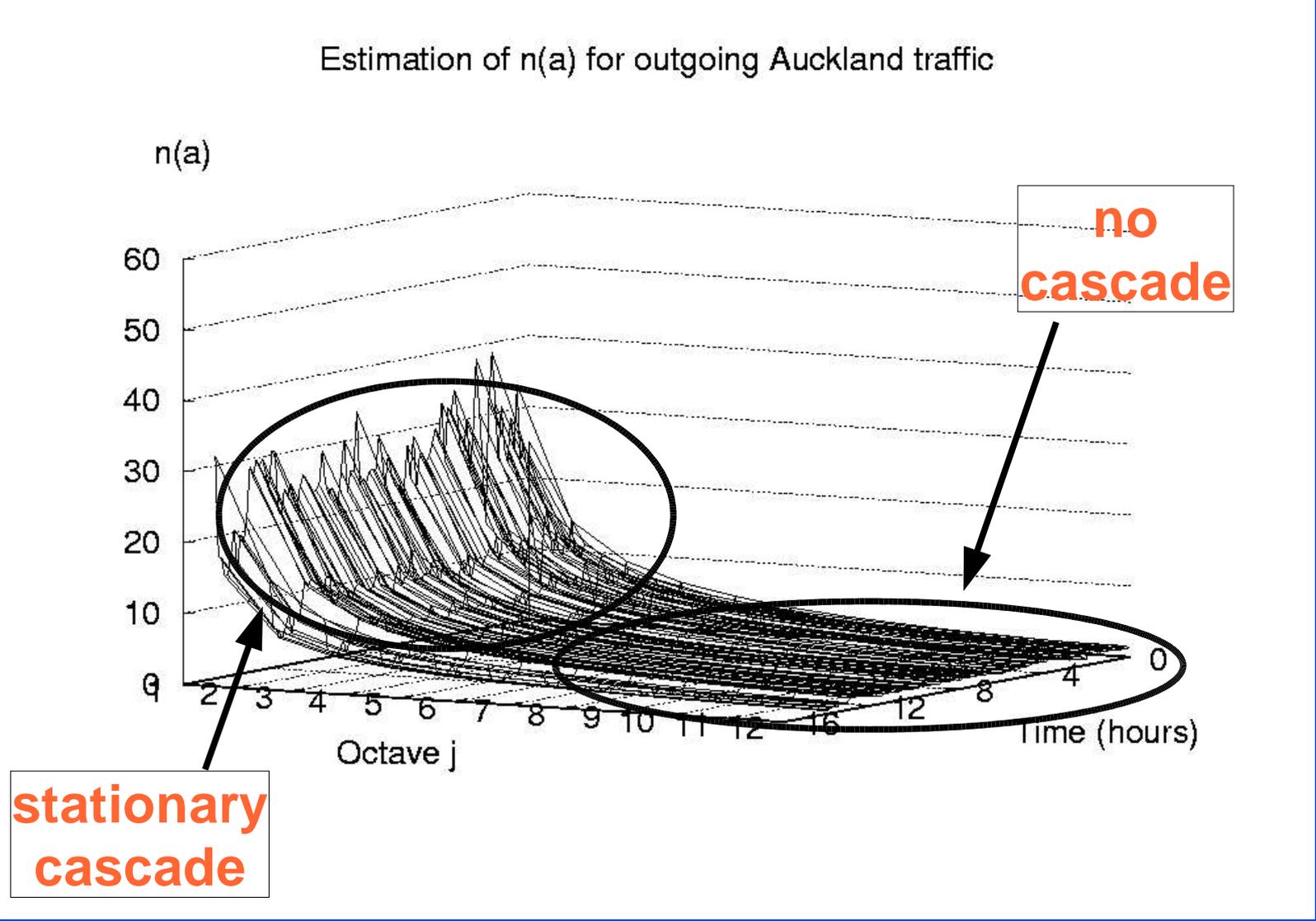
# Multiscaling analysis (in)



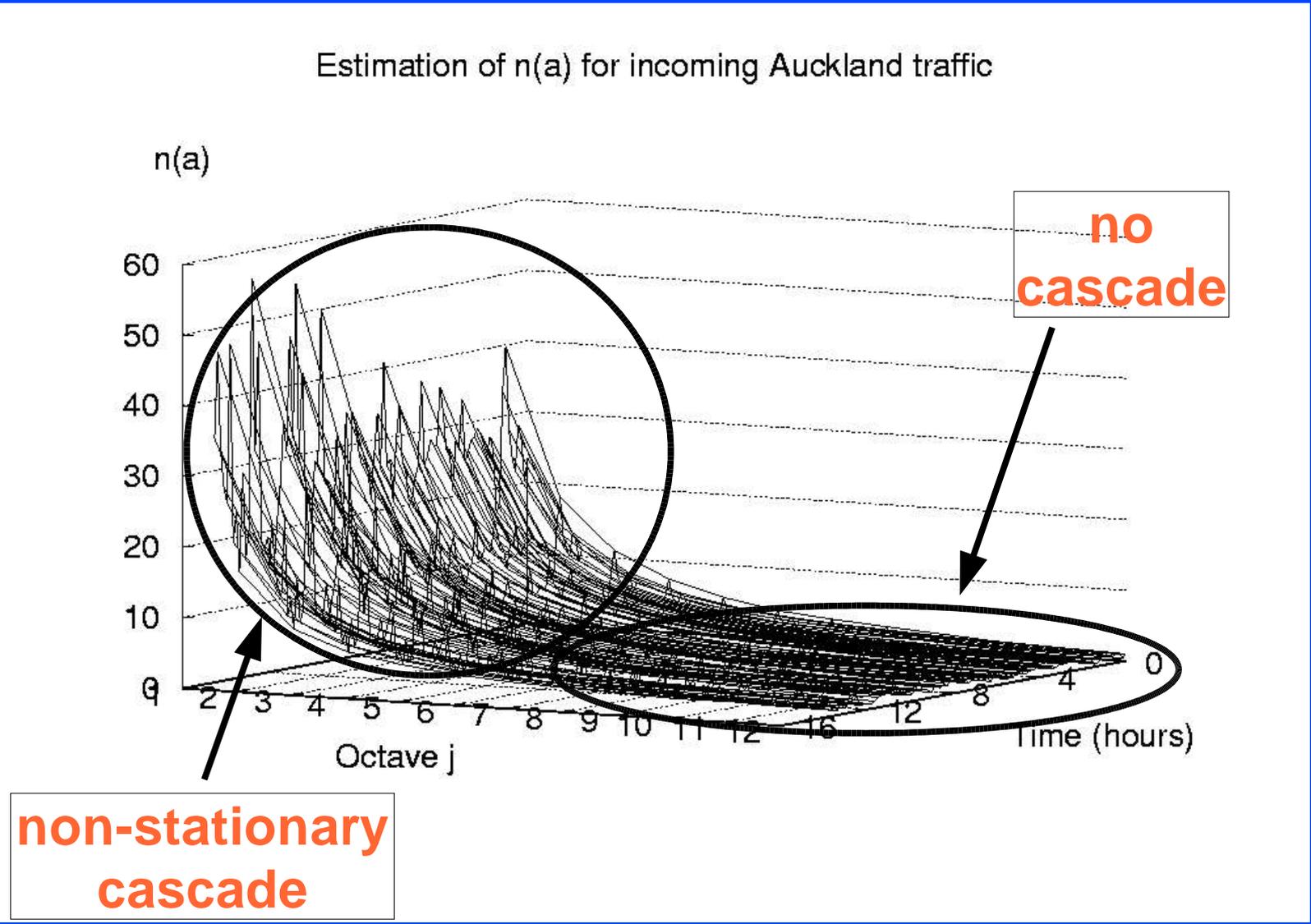
multi-scaling

mono-scaling

# Cascade analysis (out)



# Cascade analysis (in)

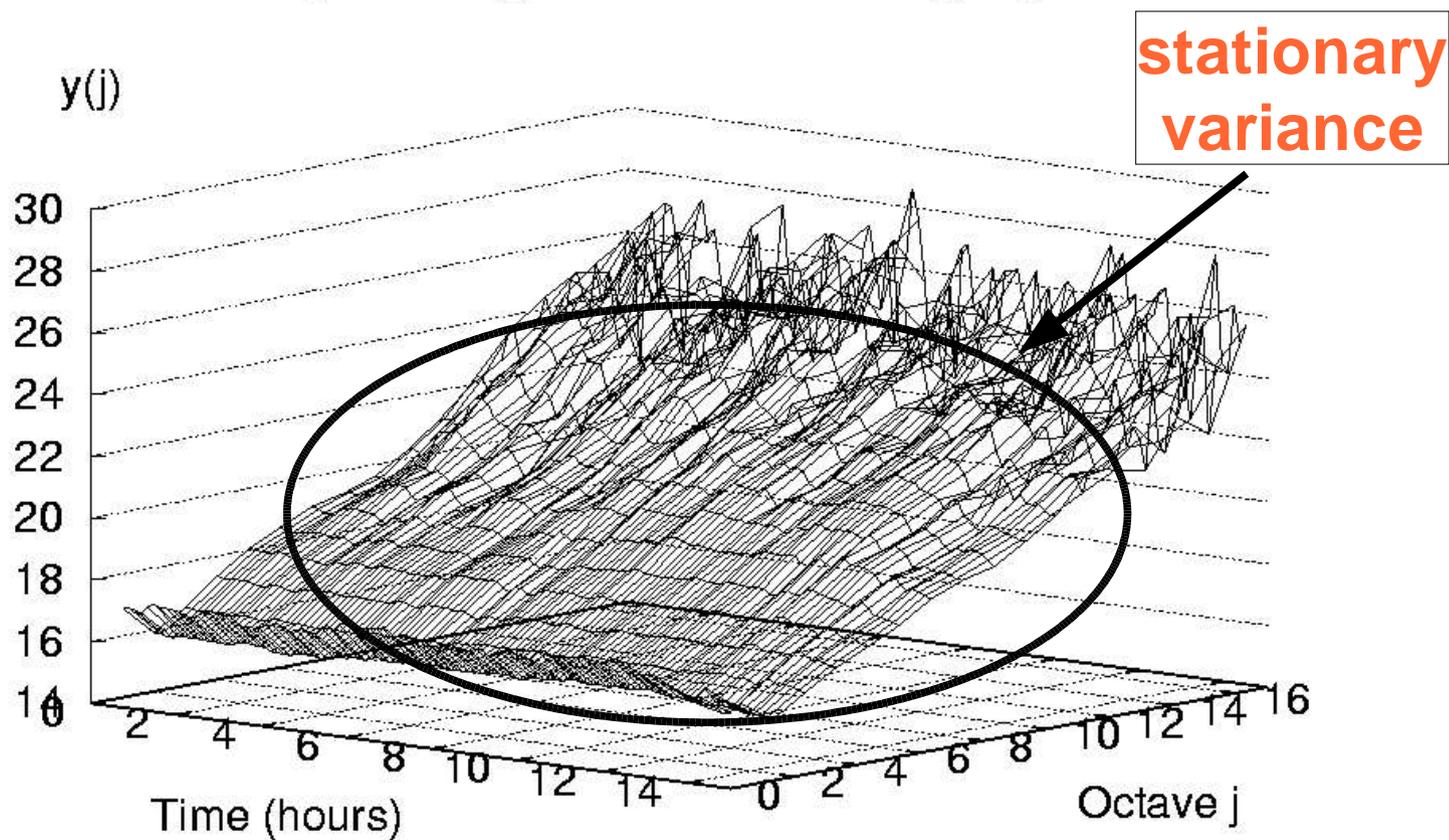


**non-stationary  
cascade**

**no  
cascade**

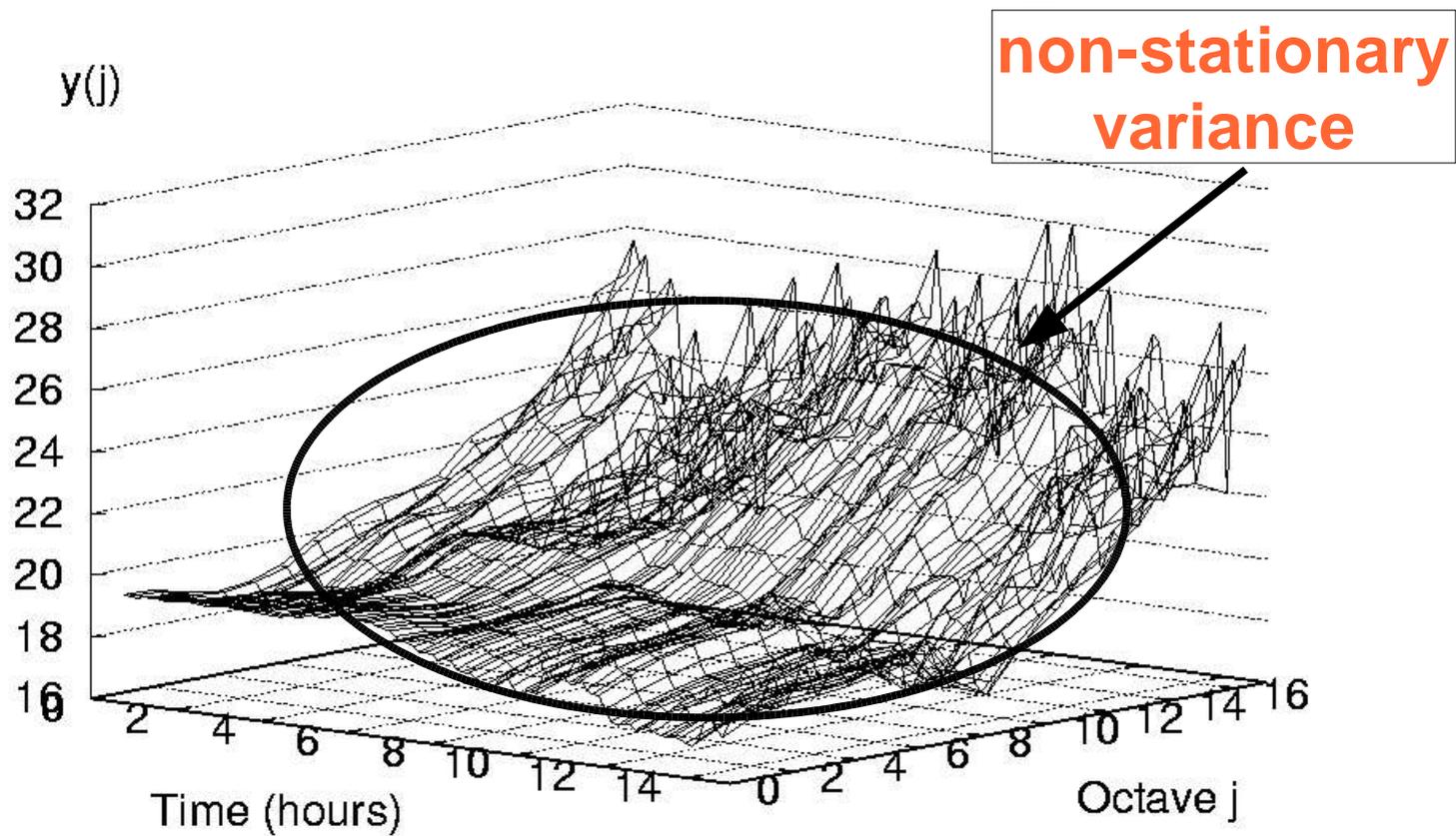
# Variance analysis (out)

3D-logscale diagram for Auckland outgoing traffic



# Variance analysis (in)

3D-logscale diagram for Auckland incoming traffic



# Conclusions

- Cascade model captures well invariance in TCP behavior
- Cascade model does not fully capture traffic dynamics, only TCP traffic segmentation
- Cascade parameters do not show effect of cross-traffic, 2<sup>nd</sup> order properties do

# After-thoughts

- Cascade is OK to generate clean TCP traffic at timescales below seconds
- Cascade is not OK to simulate realistic Internet traffic
- More work necessary to understand role of the “abstract” cascade parameters on the traffic behavior
- Can we distinguish between “normal” and “abnormal” TCP flows thanks to real-time monitoring of cascade parameters ?