

# Weighted Spectral Distribution: what a spectral topological metric can do for you?

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joint work with

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# Motivation

- Many graph measures
- Each captures specific structural property
  - local structure, e.g., clustering, communities,
  - global structure, e.g., degree distribution, assortativity
- How to compare arbitrary graphs?



# Agenda

- Graph measures
- Internet topology evolution
- Topology generators
- Application discrimination
- Dk-random graphs



# Topological measures

## Examples

- Clustering coefficient:  $\gamma(G) = \frac{1}{n} \sum_i \frac{2T_i}{d_i(d_i-1)}$ ,  $d_i \geq 2$ , with  $T_i$  the number of triangles of node  $i$
- Assortativity coefficient: correlation coefficient between node degrees
  - Internet is disassortative: large nodes connect to small ones
  - Social networks are assortative
- K-core: remaining graph when removing nodes with degree less than  $k$
- Centralities



# Weighted Spectral Distribution

## Definition

- A: adjacency matrix, D:  $\text{diag}(A)$
- Let  $\lambda_i$  be the  $i^{\text{th}}$  eigenvalue of the normalized Laplacian matrix  $L = I - D^{-1/2} A D^{1/2}$
- WSD:  $\omega(G, N) = \sum_k ((1-k)^N f(\lambda=k))$
- Depending on N, the WSD relates to specific topological properties
  - n=3: clustering
  - n=4: quasi-randomness



# Weighted Spectral Distribution

## Usage

- Graph distance metric [ToN'10]:

$$\delta(G_1, G_2) = \sum_k (1-k)^N (f_1(\lambda=k) - f_2(\lambda=k))^2$$

- WSD + random projections [Fay'10]:

$$Z = XT$$

with  $X$ : WSDs of  $m$  graphs,  $T$ : random (Gaussian) matrix

- WSD + multi-dimensional scaling [Fay'10]:
  - Build dissimilarity matrix  $R$  based on  $\delta$
  - Find approximating vectors  $Z_1, Z_2, \dots, Z_k$  that minimize a cost function  $C$



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# Internet topology evolution

- Evolving Internet structure
- Past – strong hierarchy
  - Tier-1's
  - National ISPs
  - Stubs
- Present – big players rising [Arbor'10]
  - Google
  - Akamai
  - YouTube
- What do different data-sets tell us?





# Internet topology evolution

## Inference techniques

- Control plane (e.g., BGP): routing paths towards all destinations seen by a set of observation points
- Data plane (e.g., traceroute): probe packets sent from a few boxes towards a large set of IP addresses
- Limitations:
  - BGP does not sampling all path diversity, biased by routing policies and observation points
  - Traceroute: finding a wide set of destinations to probe is hard, NAT and firewalls block probes at the edge

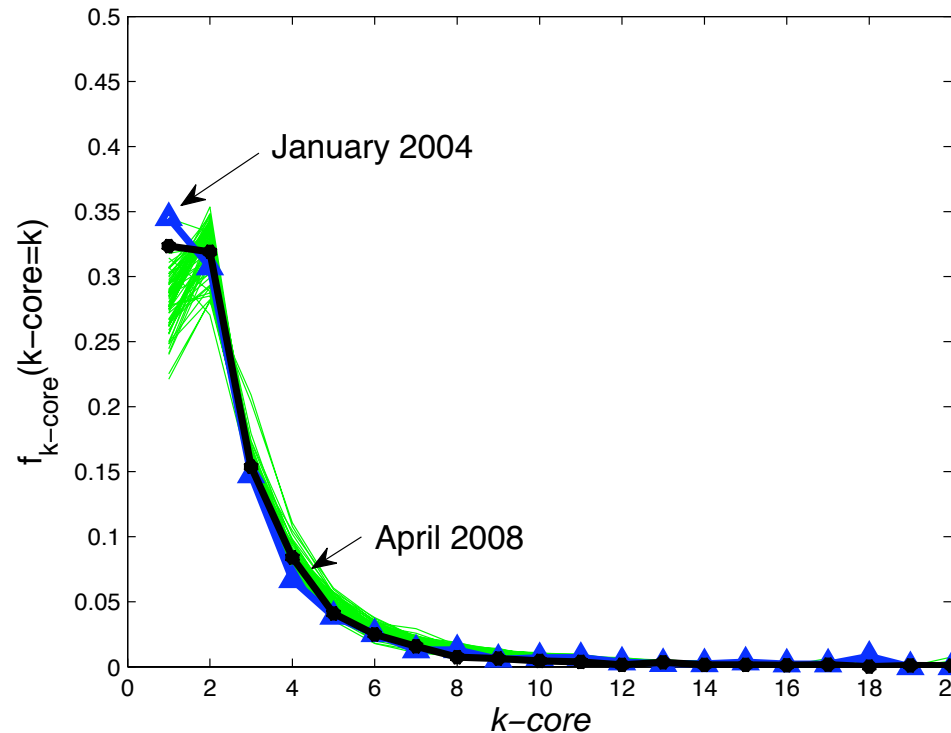
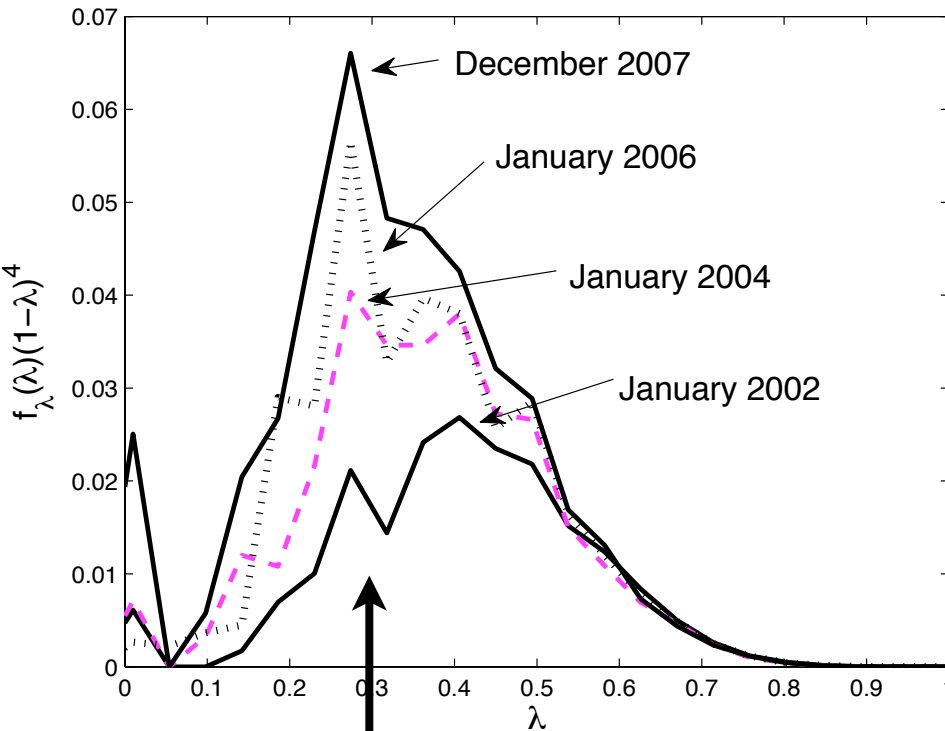


# AS topology evolution

## Traceroutes

- Reduced complexity
- More dominant core

- No significant change

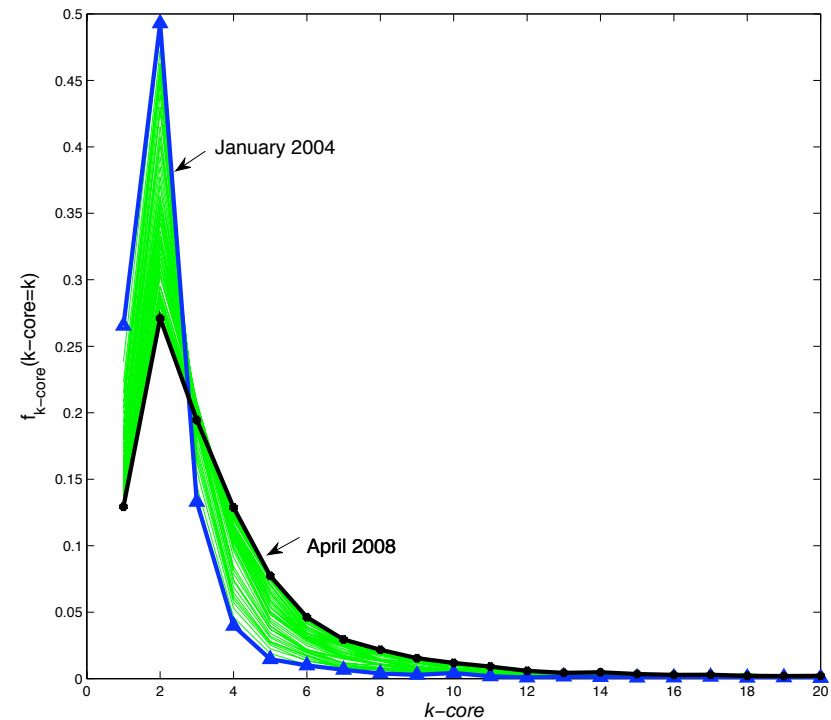
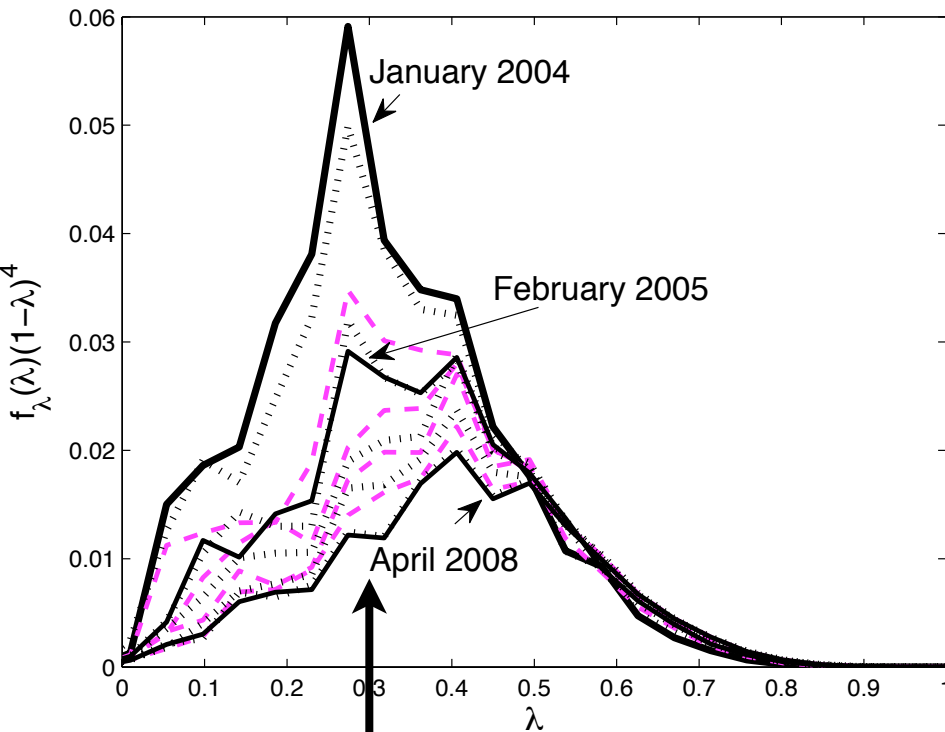


Strong hierarchy

# AS topology evolution

## BGP

- Increased complexity
- Less dominant core
- Small degree nodes get better connected



 Strong hierarchy

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- **Topology generators**
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# Generating realistic graphs

- Existing graph models
  - Random graphs
  - Power-law
  - Small world
- Each model has specific structural properties
- How randomly do existing generators sample the graph space?



# Topology generators

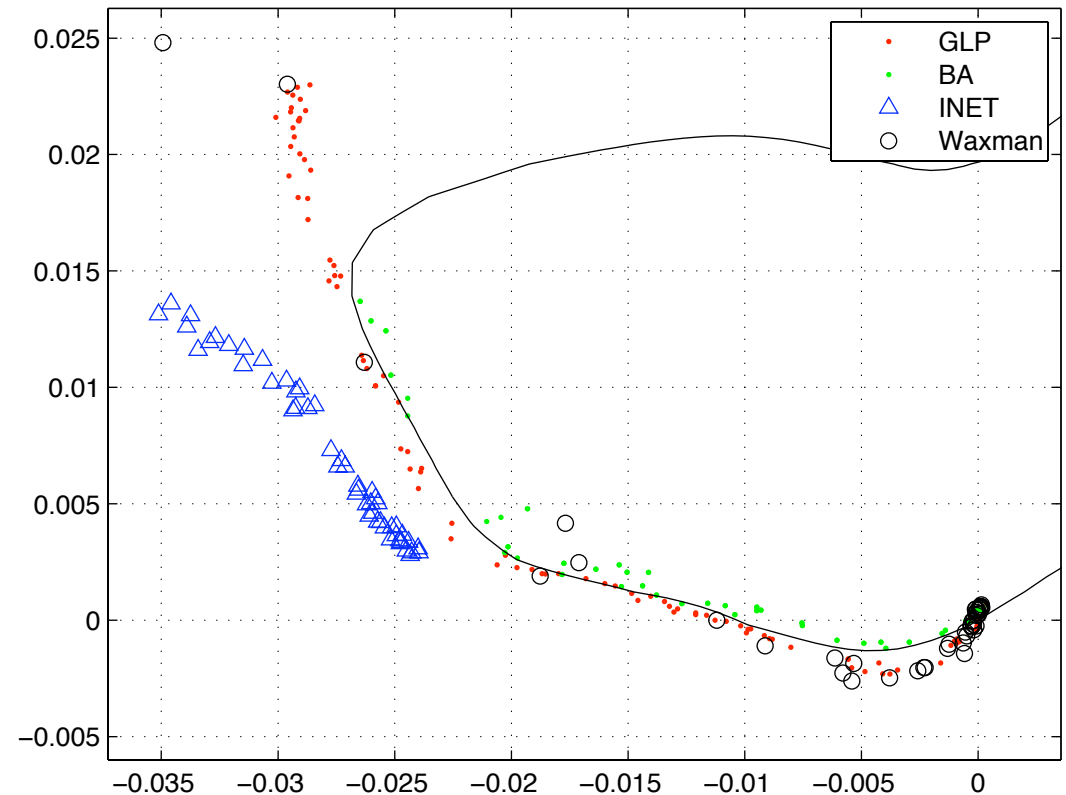
- **Waxman:** random graphs  $p(i,j) = \alpha e^{-d/(\beta L)}$   
where  $d$  is euclidian distance,  $0 < \alpha, \beta < 1$ , and  $L$  is network diameter
- **BA:** linear preferential attachment (LPA)  
 $p(i,j) = d_j / \sum_{k \in V} d_k$  where  $d_j$  is degree of node  $j$  and  $V$  is the set of nodes that have joined the network so far
- **GLP:** preferential attachment + path length and clustering match
- **Inet:** LPA + full-meshed Internet core + 30% of single-homed ASes



# Topology generators

## How random is your graph?

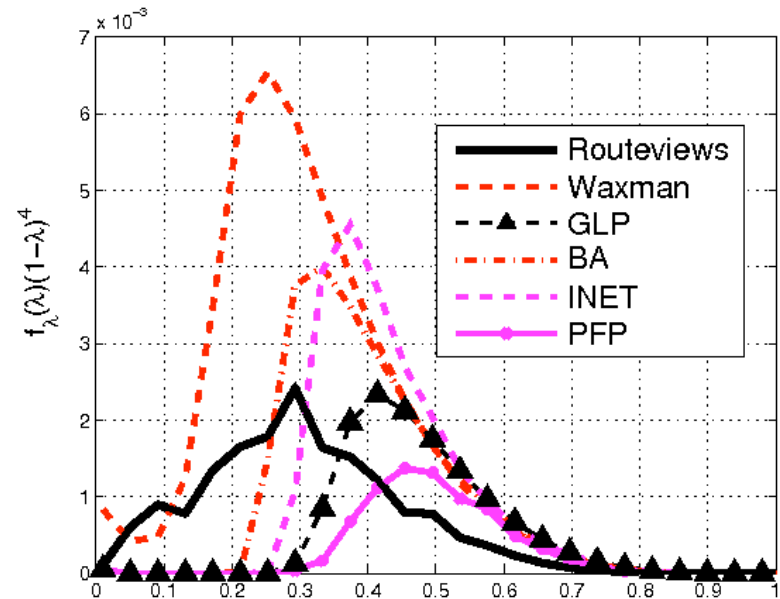
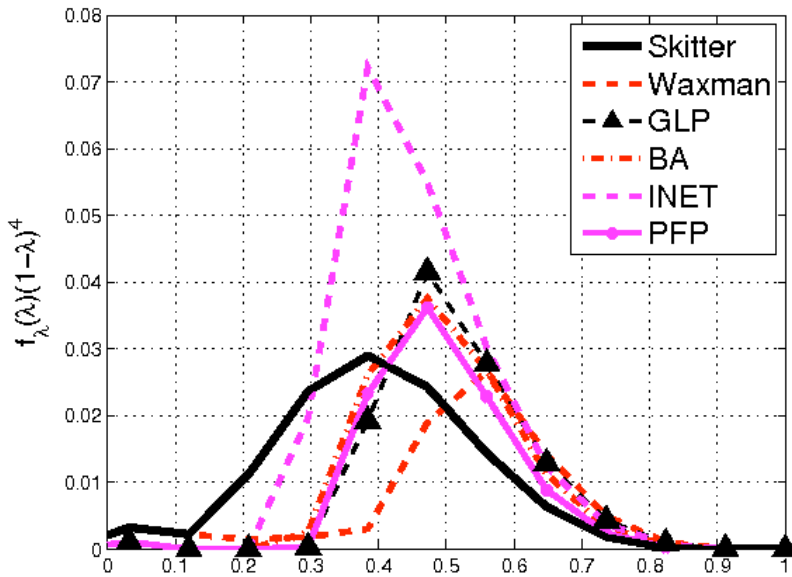
- Generate graphs with well-known topology generators
- Project their WSD onto low-dimensional space
- Each generator samples graph space in a very specific way



# Topology generators

## Optimization

- Optimize generator parameters to minimize distance with a given topology
- No generator manages to closely match any given graph





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- Topology generators
- Application discrimination
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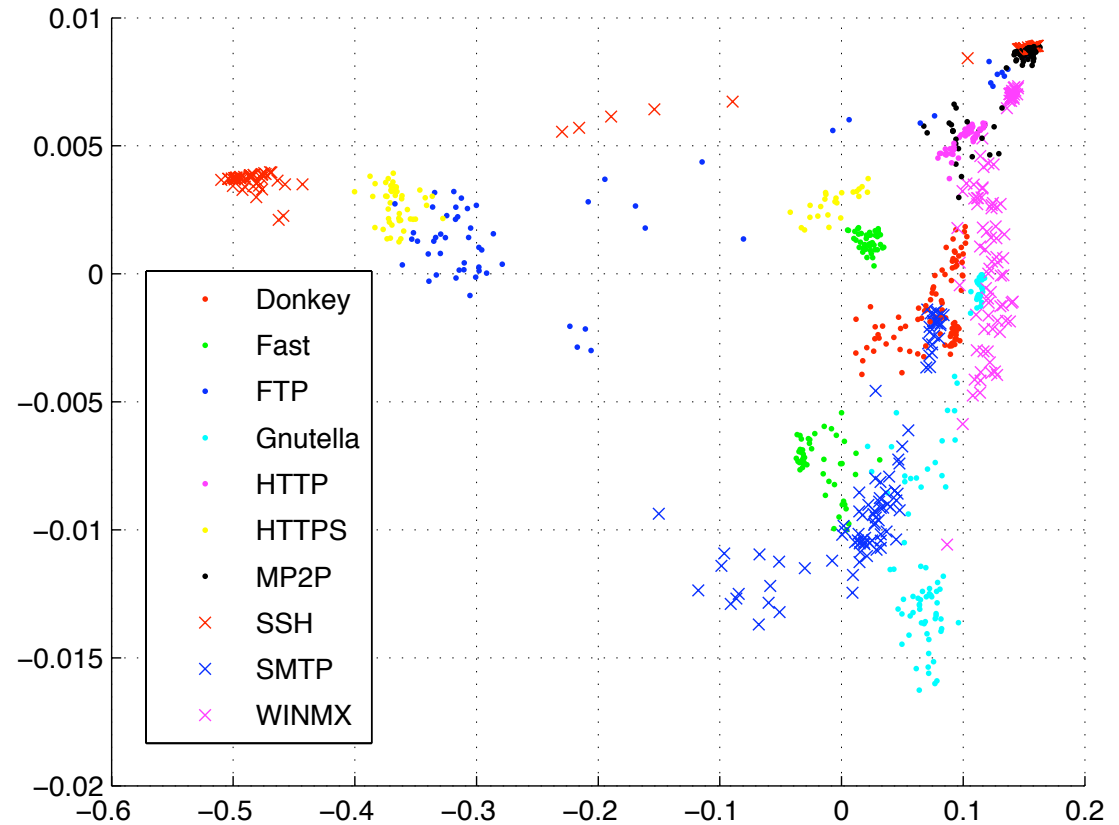
# Application identification

- Different applications have different communication patterns
- Can we discriminate between them only based on their structure?
- [CoNEXT'09] requires about 10 metrics to classify such graphs
- How many dimensions are really necessary?

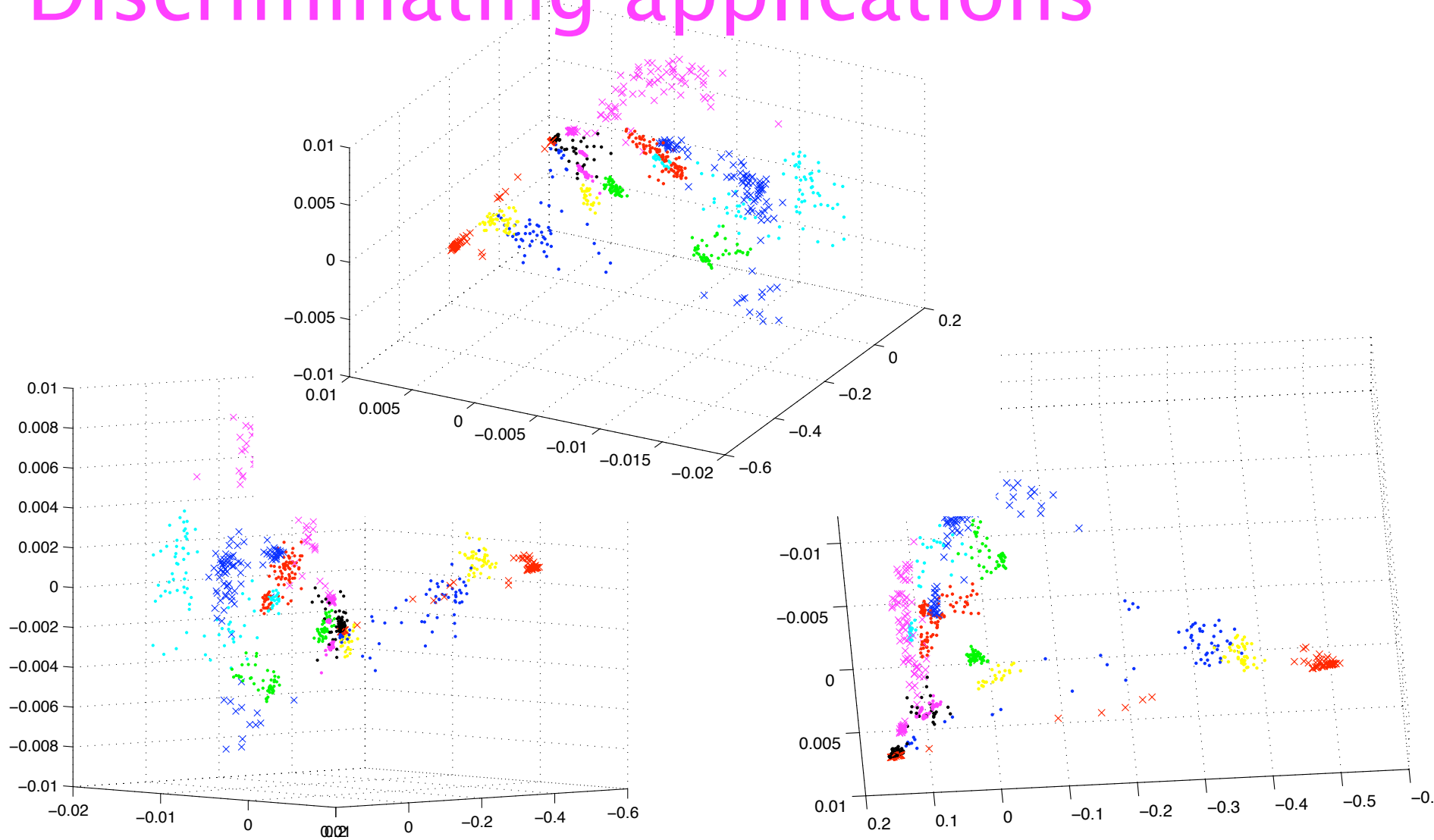


# Discriminating applications

- Using WSD and projecting on 2 dimensions already discriminates
- 86% of the flows are correctly classified
- Misclassification occurs due to low dimensionality of projection and similarity of specific applications



# Discriminating applications



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# Dk-random graphs

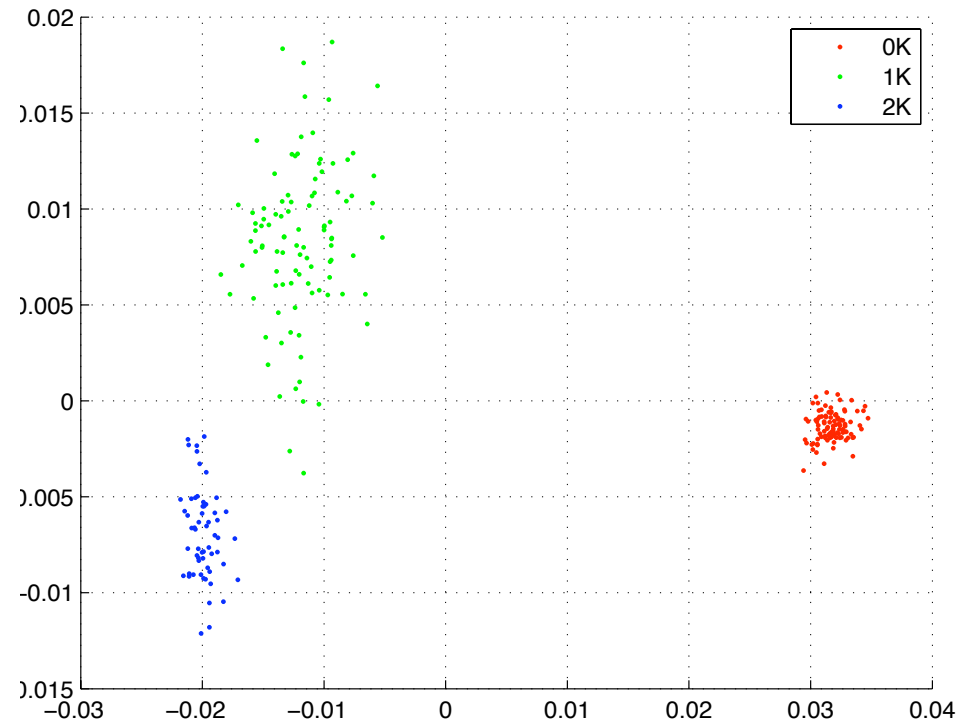
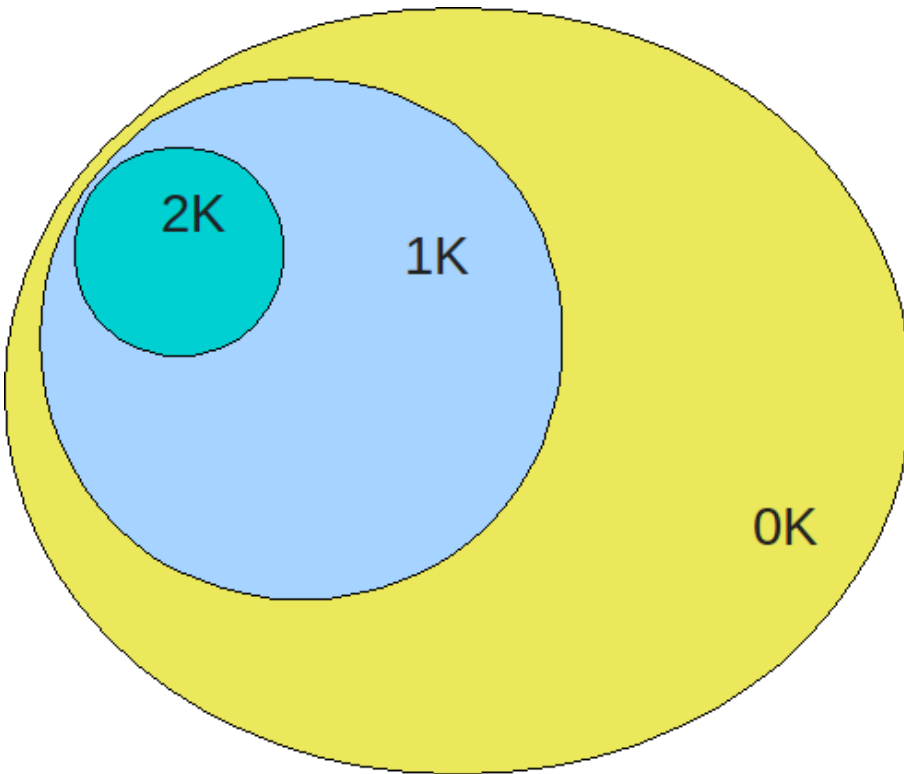
- Configuration model [Bollobas'80]
  - Builds graph with given degree distribution
- Orbis topology generator relies on dK series:
  - 0K: set of graph with average degree k
  - 1K: set of graphs with degree distribution  $p(k)$
  - 2K: set of graphs with joint degree  $p(k,j)$
  - ...
- Are graphs generated by Orbis random?



# Dk-random graphs

- Theoretical space of Dk-random graphs

- Generate 100 graphs from 0k, 1k, and 2k
- WSD + MDS



# Conclusions

- Sets of graphs are not random
  - Do not assume that topology generators sample randomly a given graph space
  - Real-world graphs even less so
- Graphs metrics like WSD can do a lot for you
  - Compare structural evolution through specific property, e.g. mixing
  - Find most appropriate topology generator
  - Discriminate specific graphs, e.g. Internet applications
  - Assess randomness in graph space, e.g. Dk-series





# Selected references

- [Arbor'10] C. Labovitz et al. Internet Inter-domain traffic. Proc. of ACM SIGCOMM, 2010.
- [Bollobas'80] B. Bollobas. A probabilistic proof of an asymptotic formula for the number of labelled regular graphs. Eur. J. Combinatorics, vol. 1, 1980.
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- [Fay'10] D. Fay et al. Discriminating graphs through spectral projections. Under submission, 2010.

