

Functional Interpretations of Intuitionistic Linear Logic

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Abstract. We present three different functional interpretations of *intuitionistic linear logic* and show how these correspond to well-known functional interpretations of intuitionistic logic via embeddings of ILL^ω into ILL^ω . The main difference from previous work of the second author is that in intuitionistic linear logic (as opposed to classical linear logic) the interpretations of $!A$ are simpler and simultaneous quantifiers are no longer needed for the characterisation of the interpretations. We then compare our approach in developing these three proof interpretations with the one of de Paiva around the Dialectica category model of linear logic.

1 Introduction

This paper presents a family of functional interpretations of *intuitionistic linear logic*. First, we present a single functional interpretation of pure (i.e., the exponential-free fragment of) intuitionistic linear logic. This is followed by a parametrised interpretation of the exponential $!A$. Finally, three possible instances of the parameter are considered and shown to correspond to three well-known functional interpretation of intuitionistic logic.

The second author [10–13] has recently shown how different functional interpretations of intuitionistic logic can be factored into a uniform family of interpretations of classical linear logic combined with Girard’s standard embedding $(\cdot)^*$ of intuitionistic logic into linear logic (see also [5]). In the symmetric context of *classical linear logic* each formula A is associated with a simultaneous one-move two-player game $|A|_y^x$. Intuitively, the two players, say Eloise and Abelard, must pick their moves x and y simultaneously and Eloise wins if and only if $|A|_y^x$ holds. The symmetric nature of the game implies that (proof-theoretically) the formula A is interpreted as the formula

$$\exists_y^x |A|_y^x$$

where $\exists_y^x A$ is a simple form of branching quantifier – termed *simultaneous quantifier* in [11]. Following this game-theoretic reading, the different interpretations of the modality $!A$ are all of the following form: First, it (always) turns a symmetric game into an asymmetric one, where Eloise plays first, giving Abelard the advantage of playing second. In the symmetric context, this asymmetric game can be modelled by allowing Abelard to play a function f which calculates his move from a given Eloise move x . Secondly, the game $!A$ gives a further (non-canonical) advantage to Abelard, by allowing him to play a *set of moves*, rather than a single move. The idea is the following: Abelard wins

the game $!A$ if there is a move $y \in fx$ that is winning with respect to Eloise’s move x , i.e. $\neg|A|_y^x$. Formally

$$|!A|_f^x \equiv \forall y \in fx |A|_y^x.$$

Therefore, the game $!A$ always introduces a break of symmetry, but it leaves open *what kind of sets* Abelard is allowed to play. What the second author has shown is that if only singleton sets are allowed the resulting interpretation corresponds to Gödel’s Dialectica interpretation [1, 7, 12]; if finite sets are allowed then it corresponds to the Diller-Nahm variant of the Dialectica interpretation [4, 13]; and if these sets are actually the whole set of moves then it corresponds to Kreisel’s modified realizability interpretation [9, 11].

In the present paper we show that in the context of *intuitionistic linear logic* every formula can be interpreted as a game where Eloise plays first and Abelard plays second, the branching quantifiers being no longer needed. In other words, Abelard’s advantage of playing second, which was limited to the game $!A$ in classical linear logic, is ubiquitous in intuitionistic linear logic. In this way, the game-theoretic interpretation of the modality $!A$ is simply to lift the moves of Abelard from a single move to a set of moves. Formally,

$$|!A|_a^x \equiv \forall y \in a |A|_y^x.$$

Therefore, by working in the context of ILL^ω , we can fully separate the canonical part of the interpretation (pure intuitionistic linear logic), where all interpretations coincide, and the non-canonical part where each choice of “sets of moves” gives rise to a different functional interpretation.

As we shall see, the functional interpretation of *pure intuitionistic linear logic* coincides with Gödel’s Dialectica interpretation of *intuitionistic logic*, reading \multimap , \otimes and \oplus as \rightarrow , \wedge and \vee , respectively. This is so because the Dialectica interpretation identifies the games A and $!A$. The connection between Gödel’s Dialectica interpretation and intuitionistic linear logic was first studied by de Paiva [14]. One can view our work here as a proof-theoretic reading of de Paiva’s category-theoretic work, together with an extension linking the “Dialectica” interpretation of intuitionistic linear logic also with Kreisel’s modified realizability (see also Biering’s recent work [2]).

The paper is organised as follows: In Section 2 we present the basic interpretation of pure intuitionistic linear logic. In the same section we outline which principles are needed for the characterisation of the interpretation (Subsection 2.1). Section 3 describes three different interpretations of the modality $!A$. This is followed (Section 4) by a description of how these choices correspond to three well-known functional interpretations of intuitionistic logic: Kreisel’s modified realizability, Diller-Nahm interpretation and Gödel’s Dialectica interpretation. Finally, in Section 5 we compare our approach (based on finite types) with that of de Paiva (based on cartesian closed categories).

1.1 Intuitionistic Linear Logic

Intuitionistic linear logic can be viewed as a fragment of Girard’s linear logic [6] which is sufficient for embedding intuitionistic logic into the linear context. We will make use of the formulation of intuitionistic linear logic shown in Tables 1 and 2 with the usual

$\frac{}{A \vdash A} \text{ (id)}$	$\frac{}{\Gamma, 0 \vdash A}$	
$\frac{\Gamma \vdash A \quad \Delta, A \vdash B}{\Gamma, \Delta \vdash B} \text{ (cut)}$	$\frac{\Gamma \vdash A}{\pi\{\Gamma\} \vdash A} \text{ (per)}$	
$\frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \otimes B} \text{ (\otimes R)}$	$\frac{\Gamma, A, B \vdash C}{\Gamma, A \otimes B \vdash C} \text{ (\otimes L)}$	
$\frac{\Gamma, A \vdash B}{\Gamma \vdash A \multimap B} \text{ (\multimap R)}$	$\frac{\Gamma \vdash A \quad \Delta, B \vdash C}{\Gamma, \Delta, A \multimap B \vdash C} \text{ (\multimap L)}$	
$\frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \& B} \text{ (& R)}$	$\frac{\Gamma, A \vdash B}{\Gamma, A \& C \vdash B} \text{ (& L)}$	$\frac{\Gamma, A \vdash B}{\Gamma, C \& A \vdash B} \text{ (& L)}$
$\frac{\Gamma \vdash A}{\Gamma \vdash A \oplus B} \text{ (\oplus R)}$	$\frac{\Gamma \vdash B}{\Gamma \vdash A \oplus B} \text{ (\oplus R)}$	$\frac{\Gamma, A \vdash C \quad \Gamma, B \vdash C}{\Gamma, A \oplus B \vdash C} \text{ (\oplus L)}$

Table 1. Intuitionistic Linear Logic (connectives)

side conditions in the rules $\forall R$ and $\exists L$. Our system is denoted by ILL^ω since we work in the language of all finite types.

The finite types are inductively defined in the usual way: i is a finite type and if ρ and σ are finite types then $\rho \rightarrow \sigma$ is a finite type. The terms of ILL^ω are: the constants (including one of type i to ensure that all types are inhabited by a closed term and the typed combinators $\Pi^{\sigma \rightarrow \tau \rightarrow \sigma}$ and $\Sigma^{(\rho \rightarrow \sigma \rightarrow \tau) \rightarrow (\rho \rightarrow \sigma) \rightarrow \rho \rightarrow \tau}$), infinitely many variables x^ρ of each finite type ρ , and if $t^{\sigma \rightarrow \tau}$ and s^σ are terms then the application ts is a term of type τ . We assume a neutral treatment of equality in the system ILL^ω (cf. [18]), i.e. the combinators are axiomatised as

$$A[\Pi xy/w] \circ\!\!\circ A[x/w] \quad \text{and} \quad A[\Sigma xyz/w] \circ\!\!\circ A[xz(yz)/w].$$

where $A \circ\!\!\circ B$ is an abbreviation from $(A \multimap B) \& (B \multimap A)$. By *combinatorial completeness*, we know that we can associate to each term t^σ and variable x^τ a term $\lambda x.t$ of type $\tau \rightarrow \sigma$ also satisfying $A[(\lambda x.t)(s)/w] \circ\!\!\circ A[t[s/x]/w]$.

The atomic formulas of ILL^ω are denoted by A_{at} (the linear logic constant 0 is an atomic formula) and if A and B are formulas, then $A \otimes B$, $A \& B$, $A \oplus B$, $A \multimap B$, $!A$, $\forall xA(x)$ and $\exists xA(x)$ are also formulas.

In this paper we will also work with a subsystem of ILL^ω , dubbed ILL_r^ω , where the following restriction is assumed on the $\&R$ -rule: The context Γ must consist entirely of formulas of the kind $!A$. In Section 4 we will see why we need this technical restriction. Nevertheless, note that both systems ILL^ω and ILL_r^ω are strong enough to capture intuitionistic logic IL^ω into the linear context, as made precise in the following proposition.

$\frac{\Gamma \vdash A}{\Gamma \vdash \forall x^p A} (\forall R)$	$\frac{\Gamma, A[t^p/x] \vdash B}{\Gamma, \forall x^p A \vdash B} (\forall L)$		
$\frac{\Gamma \vdash A[t^p/x]}{\Gamma \vdash \exists x^p A} (\exists R)$	$\frac{\Gamma, A \vdash B}{\Gamma, \exists x^p A \vdash B} (\exists L)$		
$\frac{\Gamma, !A, !A \vdash B}{\Gamma, !A \vdash B} (\text{con})$	$\frac{\Gamma \vdash B}{\Gamma, !A \vdash B} (\text{wkn})$	$\frac{! \Gamma \vdash A}{! \Gamma \vdash !A} (!R)$	$\frac{\Gamma, A \vdash B}{\Gamma, !A \vdash B} (!L)$

Table 2. Intuitionistic Linear Logic (quantifiers and modality)

Proposition 1 ([6]). Define two translations of formulas of \mathbb{L}^ω into formulas of \mathbb{ILL}_r^ω inductively as follows:

$$\begin{array}{ll}
A_{\text{at}}^* & ::= A_{\text{at}} & A_{\text{at}}^\circ & ::= !A_{\text{at}}, \text{ if } A_{\text{at}} \neq \perp \\
\perp^* & ::= 0 & \perp^\circ & ::= 0 \\
(A \wedge B)^* & ::= A^* \& B^* & (A \wedge B)^\circ & ::= A^\circ \otimes B^\circ \\
(A \vee B)^* & ::= !A^* \oplus !B^* & (A \vee B)^\circ & ::= A^\circ \oplus B^\circ \\
(A \rightarrow B)^* & ::= !A^* \multimap B^* & (A \rightarrow B)^\circ & ::= !(A^\circ \multimap B^\circ) \\
(\forall x A)^* & ::= \forall x A^* & (\forall x A)^\circ & ::= !\forall x A^\circ \\
(\exists x A)^* & ::= \exists x !A^* & (\exists x A)^\circ & ::= \exists x A^\circ
\end{array}$$

If A is provable in \mathbb{L}^ω then A^* and A° are provable in \mathbb{ILL}_r^ω (and hence also in \mathbb{ILL}^ω). Moreover, it is easy to check that $A^\circ \multimap !A^*$.

Proof. It is already known that if $\Gamma \vdash_{\mathbb{L}^\omega} A$ then $! \Gamma^* \vdash_{\mathbb{ILL}^\omega} A^*$ (see [6, 17]). The result with \mathbb{ILL}^ω replaced by \mathbb{ILL}_r^ω just requires our attention in the rule $\&R$. The result for A° follows immediately from the fact that in \mathbb{ILL}_r^ω we can prove $A^\circ \multimap !A^*$. \square

The systems \mathbb{ILL}^ω and \mathbb{ILL}_r^ω will be called *interpreted systems*, to distinguish them from the *verifying system* presented in the next subsection. The interpretations we will discuss map formulas and proofs in the interpreted system into formulas and proofs of the verifying system. In order to obtain a general notion of interpretation, we must work with the simplest (yet relevant) interpreted system possible. When extending the interpretation to more complex systems we must then ensure that the extra axioms and rules are also interpreted, but these might be interpreted by some interpretation but not others. For instance, we chose a neutral treatment of equality in the interpreted system because that is what can be interpreted in general, by all three interpretations considered. If one were to add full extensionally the Dialectica interpretation would no longer work, whereas this would be no problem for the realizability interpretation.

Notation. Throughout the paper, boldface letters $\mathbf{a}, \mathbf{b}, \dots$ or $\mathbf{x}, \mathbf{y}, \dots$ stand for tuples of terms or variables.

1.2 Verifying system

As we will show in the next sections, the three presented functional interpretations translate the formula $A \oplus B$ via a sort of flagged disjoint union, i.e. a boolean and a witness for either A or B . Therefore, in the verifying system, which we shall denote by ILL_b^ω , we assume that the language also contains the booleans b as base type. We also assume the existence of two boolean constants true and false (T, F), boolean variables, an equality relation $=^b$ between terms of boolean type, and a constant of type $b \rightarrow \rho \rightarrow \rho \rightarrow \rho$ that should be seen as a conditional λ -term $z(t, q)$ that equals either t or q depending on whether z^b equals true or false. ILL_b^ω is assumed to contain the following axioms for equality:

1. $!(x =^b x)$
2. $!(x =^b y) \multimap !(y =^b x)$
3. $!(x =^b y) \otimes !(y =^b z) \multimap !(x =^b z)$
4. $!(x =^b y) \otimes A[x/w] \multimap A[y/w]$.

We would also like to ensure that true and false are distinct and that there are no other elements of boolean type

5. $!(T =^b F) \multimap 0$
6. $!(z =^b T) \oplus !(z =^b F)$.

The axioms for the conditional λ -term are as follows

7. $A[T(t, q)/w] \multimap A[t/w]$ and $A[F(t, q)/w] \multimap A[q/w]$.

For simplicity, we use the following abbreviation:

$$A \diamond_z B := (!(z =^b T) \multimap A) \& (!(z =^b F) \multimap B).$$

Lemma 1. *The following are derivable in ILL_b^ω*

- (i)
$$\frac{\vdash A[T] \quad \vdash A[F]}{\vdash A[z]}$$
- (ii) $!(T =^b T) \multimap A \vdash A$ and $!(F =^b F) \multimap A \vdash A$
- (iii) $A \vdash !(T =^b F) \multimap B$
- (iv) $A \diamond_T B \multimap A$ and $A \diamond_F B \multimap B$
- (v) $!A \diamond_z !B \multimap !(A \diamond_z !B)$.

Proof. Assertion (i) can be derived from axioms 4. and 6.; (ii) follows easily from axiom 1.; (iii) can be deduced from axiom 5. and the forward implications in (iv) follow immediately from item (ii) and the inverse implications can easily be deduced using (iii). The forward implication in assertion (v) can be derived using assertions (i) and (iv), the other implication being trivial. \square

We stress again that we do not need to worry about which axioms are added to the *verifying system* ILL_b^ω , as these do not need to be interpreted. For instance, in the verifying system we could even have assumed full extensionality. What we listed above in the description of ILL_b^ω is the minimal necessary to verify the basic interpretation of ILL^ω and ILL_r^ω , to be described in the following section. When extending the basic interpretation to deal with $!A$ we will also need to extend the verifying system ILL_b^ω . The extensions of ILL_b^ω , however, will depend on the particular interpretation of $!A$, and will be introduced in Proposition 4 (Section 3).

2 A Basic Interpretation of Pure ILL^ω

In this section we present a basic functional interpretation of pure (without the exponential $!A$) intuitionistic linear logic, and prove its soundness. In the next section we then consider different extensions of this interpretation to full intuitionistic linear logic.

Definition 1 (Basic functional interpretation of pure ILL^ω). *For each formula A of pure ILL^ω , let us associate a formula $|A|_y^x$ of ILL_b^ω , with two fresh lists of free-variables x and y , inductively as follows: For atomic formulas A_{at} we let $|A_{\text{at}}| := A_{\text{at}}$. Assume the interpretations of A and B have already been defined as $|A|_y^x$ and $|B|_w^v$, we then define*

$$\begin{aligned} |A \multimap B|_{x,w}^{f,g} &:= |A|_{fxw}^x \multimap |B|_w^{gx} \\ |A \otimes B|_{y,w}^{x,v} &:= |A|_y^x \otimes |B|_w^v \\ |A \& B|_{y,w,z}^{x,v} &:= |A|_y^x \diamond_z |B|_w^v \\ |A \oplus B|_{y,w}^{x,v,z} &:= |A|_y^x \diamond_z |B|_w^v \\ |\exists z A(z)|_y^{x,z} &:= |A(z)|_y^x \\ |\forall z A(z)|_{y,z}^f &:= |A(z)|_y^{fz}. \end{aligned}$$

Intuitively, the meaning of A is reduced to the existence of a tuple of objects x such that $\forall y |A|_y^x$. The x 's are called *witnesses* and the y 's *challenges*. Note that, contrary to the interpretation of classical linear logic [10, 13], the functional interpretation of intuitionistic linear logic is no longer symmetric. In terms of games, the interpretation above can be seen as associating to each formula A a one-move two-player *sequential* game $|A|_y^x$. In this game, Eloise starts by playing a move x followed by Abelard playing a move y . Eloise wins if $|A|_y^x$ holds, otherwise Abelard wins.

Theorem 1 (Soundness). *Let A_0, \dots, A_n, B be formulas of pure ILL^ω , with z as the only free-variables. If*

$$A_0(\mathbf{z}), \dots, A_n(\mathbf{z}) \vdash B(\mathbf{z})$$

is provable in pure ILL^ω then terms $\mathbf{a}_0, \dots, \mathbf{a}_n, \mathbf{b}$ can be extracted from this proof such that

$$|A_0(\mathbf{z})|_{\mathbf{a}_0}^{\mathbf{x}_0}, \dots, |A_n(\mathbf{z})|_{\mathbf{a}_n}^{\mathbf{x}_n} \vdash |B(\mathbf{z})|_{\mathbf{w}}^{\mathbf{b}}$$

is provable in ILL_b^ω , with $\text{FV}(\mathbf{a}_i) \subseteq \{\mathbf{z}, \mathbf{x}_0, \dots, \mathbf{x}_n, \mathbf{w}\}$ and $\text{FV}(\mathbf{b}) \subseteq \{\mathbf{z}, \mathbf{x}_0, \dots, \mathbf{x}_n\}$.

Proof. By induction on the derivation of $A_0(\mathbf{z}), \dots, A_n(\mathbf{z}) \vdash B(\mathbf{z})$. The axioms are trivial since the interpretation does not change atomic formulas and every type is inhabited. Note that for the axiom $\Gamma, 0 \vdash A$, the interpretation of 0 is 0 itself, and so we can take arbitrary terms of the right type. The fact that every type is inhabited is also used in the study of the rules $\&\text{L}$ and $\oplus\text{R}$. The permutation rule is immediate. Let us consider a few other cases:

Cut. By induction hypothesis, assume we already terms witnessing the two premises as $|\Gamma|_{\gamma}^{\mu} \vdash |A|_{\mathbf{y}}^{\mathbf{a}_0}$ and $|\Delta|_{\delta}^{\nu}, |A|_{\mathbf{a}_1[\mathbf{x}]}^{\mathbf{x}} \vdash |B|_{\mathbf{w}}^{\mathbf{b}}$. We must construct terms that witness the conclusion $\Gamma, \Delta \vdash B$. That can be done as follows:

$$\frac{\frac{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{\mathbf{y}}^{\mathbf{a}_0}}{|\Gamma|_{\gamma'}^{\mu} \vdash |A|_{\mathbf{a}_1[\mathbf{a}_0]}^{\mathbf{a}_0}} \left[\frac{\mathbf{a}_1[\mathbf{a}_0]}{\mathbf{y}} \right]}{|\Gamma|_{\gamma'}^{\mu}, |\Delta|_{\delta'}^{\nu} \vdash |B|_{\mathbf{w}}^{\mathbf{b}'}} \frac{\frac{|\Delta|_{\delta}^{\nu}, |A|_{\mathbf{a}_1[\mathbf{x}]}^{\mathbf{x}} \vdash |B|_{\mathbf{w}}^{\mathbf{b}}}{|\Delta|_{\delta'}^{\nu}, |A|_{\mathbf{a}_1[\mathbf{a}_0]}^{\mathbf{a}_0} \vdash |B|_{\mathbf{w}}^{\mathbf{b}'}} \left[\frac{\mathbf{a}_0}{\mathbf{x}} \right]}{(\text{cut})}$$

where γ' and δ', \mathbf{b}' are obtained from γ and δ, \mathbf{b} via the substitutions $[\mathbf{a}_1[\mathbf{a}_0]/\mathbf{y}]$ and $[\mathbf{a}_0/\mathbf{x}]$, respectively.

Tensor.

$$\frac{\frac{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{\mathbf{y}}^{\mathbf{a}} \quad |\Delta|_{\delta}^{\nu} \vdash |B|_{\mathbf{w}}^{\mathbf{b}}}{|\Gamma|_{\gamma}^{\mu}, |\Delta|_{\delta}^{\nu} \vdash |A|_{\mathbf{y}}^{\mathbf{a}} \otimes |B|_{\mathbf{w}}^{\mathbf{b}}} (\otimes\text{R})}{|\Gamma|_{\gamma}^{\mu}, |\Delta|_{\delta}^{\nu} \vdash |A \otimes B|_{\mathbf{y}, \mathbf{w}}^{\mathbf{a}, \mathbf{b}}} (\text{D1})} \frac{|\Gamma|_{\gamma}^{\mu}, |A|_{\mathbf{a}}^{\mathbf{x}}, |B|_{\mathbf{b}}^{\mathbf{y}} \vdash |C|_{\mathbf{w}}^{\mathbf{c}} (\otimes\text{L})}{|\Gamma|_{\gamma}^{\mu}, |A \otimes B|_{\mathbf{a}, \mathbf{b}}^{\mathbf{x}, \mathbf{y}} \vdash |C|_{\mathbf{w}}^{\mathbf{c}}} (\text{D1})}$$

$\multimap\text{L}$ introduction.

$$\frac{\frac{\frac{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{\mathbf{y}}^{\mathbf{a}}}{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{\mathbf{y}}^{\mathbf{a}}} \left[\frac{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}{\mathbf{y}} \right]}{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}}} \quad \frac{|\Delta|_{\delta}^{\nu}, |B|_{\mathbf{b}[\mathbf{v}]}^{\mathbf{y}} \vdash |C|_{\mathbf{z}}^{\mathbf{c}[\mathbf{v}]}}{|\Delta|_{\delta}^{\nu}, |B|_{\mathbf{b}[\mathbf{g}\mathbf{a}]}^{\mathbf{g}\mathbf{a}} \vdash |C|_{\mathbf{z}}^{\mathbf{c}[\mathbf{g}\mathbf{a}]}} \left[\frac{\mathbf{g}\mathbf{a}}{\mathbf{v}} \right]}{|\Delta|_{\delta}^{\nu}, |B|_{\mathbf{b}[\mathbf{g}\mathbf{a}]}^{\mathbf{g}\mathbf{a}} \vdash |C|_{\mathbf{z}}^{\mathbf{c}[\mathbf{g}\mathbf{a}]}} (\multimap\text{L})}}{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}}, |\Delta|_{\delta}^{\nu}, |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}} \multimap |B|_{\mathbf{b}[\mathbf{g}\mathbf{a}]}^{\mathbf{g}\mathbf{a}} \vdash |C|_{\mathbf{z}}^{\mathbf{c}[\mathbf{g}\mathbf{a}]}} (\text{D1})} \frac{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}}, |\Delta|_{\delta}^{\nu}, |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}} \multimap |B|_{\mathbf{b}[\mathbf{g}\mathbf{a}]}^{\mathbf{g}\mathbf{a}} \vdash |C|_{\mathbf{z}}^{\mathbf{c}[\mathbf{g}\mathbf{a}]}}{|\Gamma|_{\gamma}^{\mu} \vdash |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}}, |\Delta|_{\delta}^{\nu}, |A|_{f\mathbf{a}(b[\mathbf{g}\mathbf{a}])}^{\mathbf{a}} \multimap |B|_{\mathbf{b}[\mathbf{g}\mathbf{a}]}^{\mathbf{g}\mathbf{a}} \vdash |C|_{\mathbf{z}}^{\mathbf{c}[\mathbf{g}\mathbf{a}]}} (\text{D1})}$$

Universal quantifier.

$$\frac{\frac{|\Gamma|_{\gamma}^{\mu} \vdash |A(\mathbf{z})|_{\mathbf{y}}^{\mathbf{a}[\mathbf{z}]}}{|\Gamma|_{\gamma}^{\mu} \vdash |A(\mathbf{z})|_{\mathbf{y}}^{\mathbf{a}[\mathbf{z}]}} \left[\frac{\lambda\mathbf{z}. \mathbf{a}[\mathbf{z}]}{\mathbf{z}} \right]}{|\Gamma|_{\gamma}^{\mu} \vdash |\forall \mathbf{z} A(\mathbf{z})|_{\mathbf{y}, \mathbf{z}}^{\lambda\mathbf{z}. \mathbf{a}[\mathbf{z}]}} (\text{D1})} \frac{\frac{|\Gamma|_{\gamma}^{\mu}, |A(\mathbf{t})|_{\mathbf{a}[\mathbf{x}]}^{\mathbf{x}} \vdash |B|_{\mathbf{w}}^{\mathbf{b}[\mathbf{x}]}}{|\Gamma|_{\gamma}^{\mu}, |A(\mathbf{t})|_{\mathbf{a}[\mathbf{f}\mathbf{t}]}^{\mathbf{f}\mathbf{t}} \vdash |B|_{\mathbf{w}}^{\mathbf{b}[\mathbf{f}\mathbf{t}]}} \left[\frac{\mathbf{f}\mathbf{t}}{\mathbf{x}} \right]}{|\Gamma|_{\gamma}^{\mu}, |\forall \mathbf{z} A(\mathbf{z})|_{\mathbf{a}[\mathbf{f}\mathbf{t}], \mathbf{t}}^{\mathbf{f}\mathbf{t}} \vdash |B|_{\mathbf{w}}^{\mathbf{b}[\mathbf{f}\mathbf{t}]}} (\text{D1})}$$

Existential quantifier.

$$\frac{|\Gamma|_{\gamma}^{\mu} \vdash |A(\mathbf{t})|_{\mathbf{y}}^{\mathbf{a}} (\text{D1})}{|\Gamma|_{\gamma}^{\mu} \vdash |\exists \mathbf{z} A(\mathbf{z})|_{\mathbf{y}}^{\mathbf{a}, \mathbf{t}}} \frac{|\Gamma|_{\gamma}^{\mu}, |A(\mathbf{z})|_{\mathbf{a}[\mathbf{z}]}^{\mathbf{x}} \vdash |B|_{\mathbf{y}}^{\mathbf{b}[\mathbf{z}]} (\text{D1})}{|\Gamma|_{\gamma}^{\mu}, |\exists \mathbf{z} A(\mathbf{z})|_{\mathbf{a}[\mathbf{z}]}^{\mathbf{x}, \mathbf{z}} \vdash |B|_{\mathbf{y}}^{\mathbf{b}[\mathbf{z}]}} (\text{D1})$$

&R introduction.

$$\begin{array}{c}
\frac{|I|_{\gamma_0}^u \vdash |A|_y^a \quad |A|_y^a \vdash |A|_y^a \diamond_T |B|_w^b}{|I|_{\gamma_0}^u \vdash |A|_y^a \diamond_T |B|_w^b} \\
\frac{\vdash |I|_{\gamma_0}^u \multimap |A|_y^a \diamond_T |B|_w^b}{\vdash |I|_{T(\gamma_0, \gamma_1)}^u \multimap |A|_y^a \diamond_T |B|_w^b} \text{ (Ax. 7)} \\
\frac{\vdash |I|_{T(\gamma_0, \gamma_1)}^u \multimap |A|_y^a \diamond_T |B|_w^b}{\vdash |I|_{z(\gamma_0, \gamma_1)}^u \multimap |A|_y^a \diamond_z |B|_w^b} \text{ (D1)} \\
\frac{\vdash |I|_{z(\gamma_0, \gamma_1)}^u \multimap |A|_y^a \diamond_z |B|_w^b}{\vdash |I|_{z(\gamma_0, \gamma_1)}^u \multimap |A \& B|_{y,w,z}^{a,b}} \text{ (D1)} \\
\frac{\vdash |I|_{z(\gamma_0, \gamma_1)}^u \multimap |A \& B|_{y,w,z}^{a,b}}{|I|_{z(\gamma_0, \gamma_1)}^u \vdash |A \& B|_{y,w,z}^{a,b}} \text{ (+)} \text{ (L1(i))}
\end{array}$$

where (+) is the dual case.

&L introduction and \oplus R introduction.

$$\begin{array}{c}
\frac{|I|_{\gamma}^u, |A|_a^x \vdash |B|_w^b}{|I|_{\gamma}^u, |A|_a^x \diamond_T |C|_c^v \vdash |B|_w^b} \text{ (L1(iv))} \\
\frac{|I|_{\gamma}^u, |A|_a^x \diamond_T |C|_c^v \vdash |B|_w^b}{|I|_{\gamma}^u, |A \& C|_{a,c,T}^{x,v} \vdash |B|_w^b} \text{ (D1)} \\
\frac{|I|_{\gamma}^u \vdash |A|_y^a \quad |A|_y^a \vdash |A|_y^a \diamond_T |B|_w^b}{|I|_{\gamma}^u \vdash |A|_y^a \diamond_T |B|_w^b} \text{ (cut)} \\
\frac{|I|_{\gamma}^u \vdash |A|_y^a \diamond_T |B|_w^b}{|I|_{\gamma}^u \vdash |A \oplus B|_{y,w}^{a,b,T}} \text{ (D1)}
\end{array}$$

The other &L and \oplus R are similar.

\oplus L introduction.

$$\begin{array}{c}
\frac{|I|_{\gamma_1}^u, |B|_b^v \vdash |C|_w^{e_2}}{|I|_{F(\gamma_0, \gamma_1)}^u, |A|_a^x \diamond_F |B|_b^v \vdash |C|_w^{F(e_1, e_2)}} \text{ (Ax. 7/L1(iv))} \\
\frac{|I|_{F(\gamma_0, \gamma_1)}^u, |A|_a^x \diamond_F |B|_b^v \vdash |C|_w^{F(e_1, e_2)}}{|I|_{z(\gamma_0, \gamma_1)}^u, |A|_a^x \diamond_z |B|_b^v \vdash |C|_w^{z(e_1, e_2)}} \text{ (L1(i))} \\
\frac{|I|_{z(\gamma_0, \gamma_1)}^u, |A|_a^x \diamond_z |B|_b^v \vdash |C|_w^{z(e_1, e_2)}}{|I|_{z(\gamma_0, \gamma_1)}^u, |A \oplus B|_{a,b}^{x,v,z} \vdash |C|_w^{z(e_1, e_2)}} \text{ (D1)}
\end{array}$$

where (+) is the dual case. The other rules are treated similarly. \square

2.1 Characterisation

As mentioned in the introduction, one of the main advantages of working in the context of intuitionistic linear logic is that we no longer need branching quantifiers. The asymmetry introduced in ILL^ω turns the symmetric games of classical linear logic into games where Eloise always plays first, so formulas A are interpreted as $\exists x \forall y |A|_y^x$.

Proposition 2. *The following principles, denoted by AC_l , MP_l , IP_l and EP (acronyms for linear versions of Axiom of Choice, Markov Principle, Independence of Premises and Extra Principle) characterise the basic interpretation presented above*

$$\begin{array}{l}
\text{AC}_l : \forall x \exists y A_\forall(y) \multimap \exists f \forall x A_\forall(fx) \\
\text{MP}_l : (\forall x A_{\text{qf}} \multimap B_{\text{qf}}) \multimap \exists x (A_{\text{qf}} \multimap B_{\text{qf}}) \\
\text{IP}_l : (A_\forall \multimap \exists y B_\forall) \multimap \exists y (A_\forall \multimap B_\forall) \\
\text{EP} : \forall x, v (A_{\text{qf}} \otimes B_{\text{qf}}) \multimap (\forall x A_{\text{qf}} \otimes \forall v B_{\text{qf}})
\end{array}$$

where A_{qf} , B_{qf} and A_{\forall} , B_{\forall} are quantifier-free formulas and purely universal formulas of ILL_b^ω respectively. It is also assumed that \mathbf{x} does not occur in B_{qf} , \mathbf{y} does not occur in A_{\forall} in the principle IP_l and \mathbf{v} does not occur in A_{qf} in the principle EP . Formally,

$$\text{ILL}_b^\omega + \text{AC}_l + \text{MP}_l + \text{IP}_l + \text{EP} \vdash A \multimap \exists \mathbf{x} \forall \mathbf{y} |A|_y^{\mathbf{x}}.$$

Moreover, assuming that \diamond is a primitive symbol, interpreted as in [11], the characterisation result still holds when bang does not occur in AC_l , MP_l , IP_l and EP and these principles are interpretable, i.e. denoting by P any instance of these principles, there are terms \mathbf{a} such that $\text{ILL}_b^\omega \vdash \forall \mathbf{y} |P|_y^{\mathbf{a}}$.

Proof. The linear equivalence can be proved by induction on the logical structure of A . Let us consider a few cases:

Tensor.

$$\begin{aligned} A \otimes B &\stackrel{(\text{IH})}{\multimap} \exists \mathbf{x} \forall \mathbf{y} |A|_y^{\mathbf{x}} \otimes \exists \mathbf{v} \forall \mathbf{w} |B|_w^{\mathbf{v}} \\ &\stackrel{(\text{EP})}{\multimap} \exists \mathbf{x}, \mathbf{v} \forall \mathbf{y}, \mathbf{w} (|A|_y^{\mathbf{x}} \otimes |B|_w^{\mathbf{v}}) \\ &\equiv \exists \mathbf{x}, \mathbf{v} \forall \mathbf{y}, \mathbf{w} |A \otimes B|_{y,w}^{\mathbf{x},\mathbf{v}}. \end{aligned}$$

With.

$$\begin{aligned} A \& B &\stackrel{(\text{IH})}{\multimap} \exists \mathbf{x} \forall \mathbf{y} |A|_y^{\mathbf{x}} \& \exists \mathbf{v} \forall \mathbf{w} |B|_w^{\mathbf{v}} \\ &\multimap \forall z (\exists \mathbf{x} \forall \mathbf{y} |A|_y^{\mathbf{x}} \diamond_z \exists \mathbf{v} \forall \mathbf{w} |B|_w^{\mathbf{v}}) \\ &\multimap \forall z \exists \mathbf{x}, \mathbf{v} (\forall \mathbf{y} |A|_y^{\mathbf{x}} \diamond_z \forall \mathbf{w} |B|_w^{\mathbf{v}}) \\ &\multimap \forall z \exists \mathbf{x}, \mathbf{v} \forall \mathbf{y}, \mathbf{w} (|A|_y^{\mathbf{x}} \diamond_z |B|_w^{\mathbf{v}}) \\ &\stackrel{(\text{AC}_l)}{\multimap} \exists \mathbf{f}, \mathbf{g} \forall z, \mathbf{y}, \mathbf{w} (|A|_y^{\mathbf{f}z} \diamond_z |B|_w^{\mathbf{g}z}) \\ &\multimap \exists \mathbf{x}, \mathbf{v} \forall z, \mathbf{y}, \mathbf{w} (|A|_y^{\mathbf{x}} \diamond_z |B|_w^{\mathbf{v}}) \\ &\equiv \exists \mathbf{x}, \mathbf{v} \forall \mathbf{y}, \mathbf{w}, z |A \& B|_{y,w,z}^{\mathbf{x},\mathbf{v}}. \end{aligned}$$

Linear implication.

$$\begin{aligned} A \multimap B &\stackrel{(\text{IH})}{\multimap} \exists \mathbf{x} \forall \mathbf{y} |A|_y^{\mathbf{x}} \multimap \exists \mathbf{v} \forall \mathbf{w} |B|_w^{\mathbf{v}} \stackrel{(\text{IP}_l, \text{MP}_l)}{\multimap} \forall \mathbf{x} \exists \mathbf{v} \forall \mathbf{w} \exists \mathbf{y} (|A|_y^{\mathbf{x}} \multimap |B|_w^{\mathbf{v}}) \\ &\stackrel{(\text{AC}_l)}{\multimap} \exists \mathbf{f}, \mathbf{g} \forall \mathbf{x}, \mathbf{w} (|A|_{f,x}^{\mathbf{x}} \multimap |B|_w^{\mathbf{g}x}) \equiv \exists \mathbf{f}, \mathbf{g} \forall \mathbf{x}, \mathbf{w} |A \multimap B|_{x,w}^{\mathbf{f},\mathbf{g}}. \end{aligned}$$

Universal quantifier.

$$\forall z A \stackrel{(\text{IH})}{\multimap} \forall z \exists \mathbf{x} \forall \mathbf{y} |A|_y^{\mathbf{x}} \stackrel{(\text{AC}_l)}{\multimap} \exists \mathbf{f} \forall \mathbf{y}, z |A|_y^{\mathbf{f}z} \equiv \exists \mathbf{f} \forall \mathbf{y}, z | \forall z A |_{y,z}^{\mathbf{f}}.$$

The other cases are treated similarly. In fact, for the remaining cases (once the induction hypothesis is assumed) the equivalence can be proved in ILL_b^ω alone.

With the assumptions presented, the interpretability of the principles is easily checked since quantifier-free formulas are interpretable by themselves, i.e. they do not ask for realisers. We illustrate with the principle AC_l where the premise is interpreted as

$$|\forall \mathbf{x} \exists \mathbf{y} \forall z A_{\text{qf}}|_{z,x}^{\mathbf{f}} \equiv |\exists \mathbf{y} \forall z A_{\text{qf}}(\mathbf{x}, \mathbf{y}, z)|_z^{\mathbf{f}x} \equiv |\forall z A_{\text{qf}}(\mathbf{x}, \mathbf{f}x, z)|_z \equiv A_{\text{qf}}(\mathbf{x}, \mathbf{f}x, z)$$

whereas the conclusion is interpreted as

$$|\exists f \forall x \forall z A_{qf}(x, fx, z)|_{z,x}^f \equiv |\forall x \forall z A_{qf}(x, fx, z)|_{z,x} \equiv |\forall z A_{qf}(x, fx, z)|_z \equiv A_{qf}(x, fx, z).$$

Since the realisers of the premise are the same as those of the conclusion, the identity and projection functions can be taken as realisers of the implication. \square

Remark 1. Note that if we are embedding ILL^ω via the standard embedding $(\cdot)^*$ then the connective \otimes is not needed, and hence the extra principle EP is not needed either.

3 Some Interpretations of ILL^ω

In this section we consider a few choices of how the basic interpretation given in Definition 1 can be extended to full intuitionistic linear logic, i.e. we present three possible interpretations of $!A$. All choices considered will have the form

$$|!A|_y^x := !\forall y' \sqsubset y |A|_{y'}^x, \quad (1)$$

where $\forall y \sqsubset a A$ is a meta-level formula construction which we will assume to satisfy the following: For some terms¹ $\eta(\cdot)$, $(\cdot) \otimes (\cdot)$ and $(\cdot) \circ (\cdot)$ the conditions below are provable in ILL_b^ω

- (A1) $!\forall y \sqsubset \eta(z) A[y] \multimap A[z]$
- (A2) $!\forall y \sqsubset (y_1 \otimes y_2) A[y] \multimap !(\forall y \sqsubset y_1 A[y]) \otimes !(\forall y \sqsubset y_2 A[y])$
- (A3) $!\forall y \sqsubset (f \circ z) A[y] \multimap !\forall x \sqsubset z !\forall y \sqsubset fx A[y]$.

The three instances of such meta-level formula construction $\forall y \sqsubset a A$ we will consider are $\forall y A$, $\forall y \in a A$ (where $y \in a$ will be defined later), and $A[a/y]$.

Proposition 3. *Under the assumptions (A1 – A3) on the formula construction $\forall y \sqsubset a A$, the generic interpretation of $!A$ as above leads to a sound functional interpretation of ILL^ω .*

Proof. By Theorem 1 we just have to analyse the rules of contraction, weakening, $!R$, and $!L$.

Contraction. Assume by induction hypothesis that we already have terms witnessing the premise of the rule, i.e. $|\Gamma|_\gamma^\mu, |!A|_{a_0}^{x_0}, |!A|_{a_1}^{x_1} \vdash |B|_w^b$. We must from these construct witnesses for the conclusion $\Gamma, !A \vdash B$. That can be done as follows:

$$\frac{\frac{|\Gamma|_\gamma^\mu, |!A|_{a_0}^{x_0}, |!A|_{a_1}^{x_1} \vdash |B|_w^b}{|\Gamma|_\gamma^\mu, |!A|_{a_0}^{x_0}, |!A|_{a_1}^{x_1} \vdash |B|_w^b} \left[\frac{x}{x_0}, \frac{x}{x_1} \right] (1)}{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 |A|_{y'}^x, !\forall y' \sqsubset a_1 |A|_{y'}^x \vdash |B|_w^b} (1)}{\frac{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 |A|_{y'}^x \otimes !\forall y' \sqsubset a_1 |A|_{y'}^x \vdash |B|_w^b}{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 |A|_{y'}^x \otimes !\forall y' \sqsubset a_1 |A|_{y'}^x \vdash |B|_w^b} (\otimes L)}{\frac{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 \otimes a_1 |A|_{y'}^x \vdash |B|_w^b}{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 \otimes a_1 |A|_{y'}^x \vdash |B|_w^b} (A2)}{\frac{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 \otimes a_1 |A|_{y'}^x \vdash |B|_w^b}{|\Gamma|_\gamma^\mu, !\forall y' \sqsubset a_0 \otimes a_1 |A|_{y'}^x \vdash |B|_w^b} (1)} (1)$$

¹ Note that these terms are allowed to be specific to the formula A , in particular, the free variables of $\eta(\cdot)$, $(\cdot) \otimes (\cdot)$ and $(\cdot) \circ (\cdot)$ are assumed to be contained in the free-variables of $\forall y A[y]$ (i.e. all free-variables of A except y).

Weakening.

$$\frac{\frac{|Γ|_{\gamma}^a \vdash |B|_w^b}{|Γ|_{\gamma}^a, !\forall y' \sqsubset a |A|_{y'}^x \vdash |B|_w^b} \text{ (wkn)}}{|Γ|_{\gamma}^a, !A|_a^x \vdash |B|_w^b} \text{ (1)}$$

where a are arbitrary closed terms of the appropriate types. Note that every type is inhabited by a closed term.

!R.

$$\frac{\frac{|!Γ|_{\gamma[y']}^a \vdash |A|_{y'}^a}{!\forall w' \sqsubset \gamma[y'] |Γ|_{w'}^a \vdash |A|_{y'}^a} \text{ (1)}}{!\forall y' \sqsubset y !\forall w' \sqsubset (\lambda y'. \gamma[y']) y' |Γ|_{w'}^a \vdash !\forall y' \sqsubset y |A|_{y'}^a} \text{ (A3)}$$

$$\frac{!\forall w' \sqsubset (\lambda y'. \gamma[y']) \circ y |Γ|_{w'}^a \vdash !\forall y' \sqsubset y |A|_{y'}^a}{!\Gamma|_{(\lambda y'. \gamma[y']) \circ y}^a \vdash !A|_y^a} \text{ (1)}$$

!L.

$$\frac{\frac{|Γ|_{\gamma}^a, |A|_a^x \vdash |B|_w^b}{|Γ|_{\gamma}^a, !\forall y \sqsubset \eta(a) |A|_y^x \vdash |B|_w^b} \text{ (A1)}}{|Γ|_{\gamma}^a, !A|_{\eta(a)}^x \vdash |B|_w^b} \text{ (1)}$$

That concludes the proof. \square

Remark 2. Assume that the types of y^ρ and $a^{T\rho}$ in $\forall y \sqsubset a |A|_y^x$ are as shown, for a fixed A . Then, our three families of terms have types

$$\begin{aligned} \eta &: \rho \rightarrow T\rho \\ \otimes &: T\rho \times T\rho \rightarrow T\rho \\ \circ &: (\tau \rightarrow T\rho) \times T\tau \rightarrow T\rho. \end{aligned}$$

In category theory, one could think of (T, η, \circ) as forming a Kleisli triple (\sim monad), with \otimes being a commutative monoid on $T\rho$. This in turn extends to a comonad on formulas as

$$T(A[y]) := !\forall y \sqsubset a A,$$

where the formula A with free-variables y is transformed in the new formula $!\forall y \sqsubset a A$ with free-variables a . See e.g. the work of Valeria de Paiva [15] and Martin Hyland ([8], section 3.1) on categorical logic for more information about the connection between functional interpretations and comonads. More on the relation between ours and de Paiva's work can be found in Section 5.

Next, we present three sound interpretations of $!A$ by providing three instances of $\forall y \sqsubset a A$ which satisfy conditions (A1), (A2), and (A3). It is important to observe that the meta-level formula construction $\forall y \sqsubset a A$ is part of the *verifying system*. Therefore, when discussing particular instances of $\forall y \sqsubset a A$ it is not relevant for the interpretation how the terms needed are axiomatised. Only axioms in the *interpreted system* needed to be interpreted.

Proposition 4. *We have the following:*

(a) $!A|_y^x := !\forall y|A|_y^x$ is a sound interpretation of $!A$.

(b) Assume that the language of the verifying system ILL_b^ω has a new finite type σ^* for each finite type σ . An element of type σ^* is a finite set of elements of type σ . The extended language has a relation symbol \in infixing between a term of type σ and a term of type σ^* with axioms to ensure that $!(x \in y)$ if and only if x is an element in the set y . Let then the formula $\forall x \in t A$ abbreviate $\forall x(!(x \in t) \rightarrow A)$. Assume also the existence of three more constants $\eta : \sigma \rightarrow \sigma^*$, $\otimes : \sigma^* \rightarrow \sigma^* \rightarrow \sigma^*$ and $\circ : \sigma^* \rightarrow (\sigma \rightarrow \rho^*) \rightarrow \rho^*$ that should be seen as terms such that $\eta(t)$ is the singleton set with t^σ as the only element (in particular $!(t \in \eta(t))$), $t \otimes q$ is the union of two finite sets t and q , and $f \circ q$ is the set that results from the union of all sets fx with $x \in q$. Then $!A|_a^x := !\forall y \in a |A|_y^x$ is a sound interpretation of $!A$.

(c) Assume the verifying system ILL_b^ω has an extra axiom schema $\vdash !A \oplus (!A \rightarrow 0)$, asserting the decidability of quantifier-free-formulas A . Assume also that definition by cases is definable over quantifier-free formulas A in the term language of ILL_b^ω , i.e.

$$t \otimes s := \begin{cases} t & \text{if } !A \rightarrow 0 \\ s & \text{if } !A, \end{cases}$$

with the rules

$$\frac{\Gamma \vdash B[t \otimes s]}{\Gamma, !A \vdash B[s]} \quad \frac{\Gamma \vdash B[t \otimes s]}{\Gamma, !A \rightarrow 0 \vdash B[t]}$$

Then, $!A|_y^x := !A|_y^x$ is a sound interpretation of $!A$.

Proof. (a) This interpretation of $!A$ corresponds to the choice $\forall y \sqsubset t A := \forall y A$. It is easy to check that conditions (A1), (A2) and (A3) become

$$\begin{aligned} & !\forall y A[y] \rightarrow A[z] \\ & !\forall y A[y] \rightarrow !\forall y A[y] \otimes !\forall y A[y] \\ & !\forall y A[y] \rightarrow !\forall x !\forall y A[y] \end{aligned}$$

respectively, which are trivially derivable in ILL_b^ω .

(b) The interpretation $!A|_a^x := !\forall y \in a |A|_y^x$ corresponds to the choice $\forall y \sqsubset t A := \forall y \in t A$, i.e. $\forall y(!(y \in t) \rightarrow A[y])$. In this context, the conditions (A1), (A2) and (A3) become

$$\begin{aligned} & !\forall y \in \eta(z) A[y] \rightarrow A[z] \\ & !\forall y \in y_1 \otimes y_2 A[y] \rightarrow !\forall y \in y_1 A[y] \otimes !\forall y \in y_2 A[y] \\ & !\forall y \in f \circ z A[y] \rightarrow !\forall x \in z !\forall y \in fx A[y], \end{aligned}$$

which are derivable in the extension of ILL_b^ω outlined above.

(c) This interpretation of $!A$ corresponds to the choice $\forall y \sqsubset t A[y] := A[t/y]$. Given a formula $A[y]$ we define $\eta(\cdot)$, as being the identity, \circ is defined as $f \circ x := fx$ and $y_1 \otimes y_2$ as

$$y_1 \otimes y_2 := \begin{cases} y_1 & \text{if } !A[y_1] \multimap 0 \\ y_2 & \text{if } !A[y_1]. \end{cases}$$

Conditions (A1), (A2) and (A3) then become

$$\begin{aligned} & !A[\eta(z)] \multimap A[z] \\ & !A[y_1 \otimes y_2] \multimap !A[y_1] \otimes !A[y_2] \\ & !A[f \circ z] \multimap !!A[fz] \end{aligned}$$

respectively. From the definitions of $\eta(\cdot)$ and $(\cdot) \circ (\cdot)$ conditions (A1) and (A3) are trivially derivable. In the derivation of (A2) we use

$$\begin{aligned} & \vdash !A \oplus (!A \multimap 0) \\ & !A[y_1], !A[y_1 \otimes y_2] \vdash !A[y_1] \otimes !A[y_2], \text{ and} \\ & !A[y_1] \multimap 0, !A[y_1 \otimes y_2] \vdash 0. \end{aligned}$$

More precisely,

$$\begin{array}{c} \frac{\frac{!A[y_1], !A[y_1 \otimes y_2] \vdash !A[y_1] \otimes !A[y_2] \quad !A[y_1] \multimap 0, !A[y_1 \otimes y_2] \vdash !A[y_1] \otimes !A[y_2]}{!A[y_1] \oplus (!A[y_1] \multimap 0), !A[y_1 \otimes y_2] \vdash !A[y_1] \otimes !A[y_2]} \text{ (cut)}}{!A[y_1 \otimes y_2] \vdash !A[y_1] \otimes !A[y_2]} \text{ (+)} \end{array}$$

where (+) is an instance of the assumed axiom $!A[y_1] \oplus (!A[y_1] \multimap 0)$. \square

4 Relation to Standard Interpretations of IL^ω

We argued in the introduction (see Proposition 1) that for the purpose of analysing IL^ω via linear logic it suffices to work with the system ILL_r^ω . As it turns out, in ILL_r^ω , we can simplify the interpretation of the connective $\&$, so that we no longer need the boolean variable z in \diamond_z in that particular case.

Proposition 5. *When interpreting the subsystem ILL_r^ω , the interpretation of A & B presented in Definition 1 can be simplified so that the parametrised interpretation*

$$\begin{aligned} |A \multimap B|_{x,w}^{f,g} &::= |A|_{fxw}^x \multimap |B|_w^{gx} \\ |A \otimes B|_{y,w}^{x,v} &::= |A|_y^x \otimes |B|_w^v \\ |A \& B|_{y,w}^{x,v} &::= |A|_y^x \& |B|_w^v \\ |A \oplus B|_{y,w}^{x,v,z} &::= |A|_y^x \diamond_z |B|_w^v \\ |\exists z A(z)|_y^{x,z} &::= |A(z)|_y^x \\ |\forall z A(z)|_{y,z}^f &::= |A(z)|_y^{fz} \\ |!A|_y^x &::= !\forall y' \sqsubset y |A|_{y'}^x \end{aligned}$$

is sound for ILL_r^ω , assuming (A1), (A2), and (A3) are satisfied.

Proof. We just have to analyse the rules for $\&$ having in mind that, in the case of the system under interpretation, the $\&R$ introduction is restricted of the form $! \Gamma$. The simplified interpretation of $A \& B$ is shown sound as:

$$\frac{\frac{\frac{! \Gamma_{\gamma_0}^u \vdash |A|_x^a}{! \forall y' \sqsubset \gamma_0 | \Gamma_{y'}^u \vdash |A|_x^a} \text{(P5)} \quad \frac{\frac{! \Gamma_{\gamma_1}^u \vdash |B|_y^b}{! \forall y' \sqsubset \gamma_1 | \Gamma_{y'}^u \vdash |B|_y^b} \text{(P5)}}{! \forall y' \sqsubset (\gamma_0 \otimes \gamma_1) | \Gamma_{y'}^u \vdash |B|_y^b} \text{(A2)}}{! \forall y' \sqsubset (\gamma_0 \otimes \gamma_1) | \Gamma_{y'}^u \vdash |A|_x^a \& |B|_y^b} \text{(&R)}}{! \Gamma_{\gamma_0 \otimes \gamma_1}^u \vdash |A \& B|_{x,y}^{a,b}} \text{(P5)}$$

And for the left introduction:

$$\frac{\frac{| \Gamma_{\gamma}^u, |A|_a^x \vdash |C|_w^c}{| \Gamma_{\gamma}^u, |A|_a^x \& |B|_b^v \vdash |C|_w^c} \text{(&L)}}{| \Gamma_{\gamma}^u, |A \& B|_{a,b}^{x,v} \vdash |C|_w^c} \text{(P5)}$$

The other $\&L$ introduction is similar. □

Since in the remaining part of this section we work with translations of intuitionistic logic into linear logic, by $|A|_y^x$ we refer to the (simplified) parametrised interpretation described in Proposition 5. Next we prove that the three different ways of interpreting $!A$ (cf. Proposition 4) give rise to interpretations of ILL_r^ω that correspond (via the translations of intuitionistic logic into intuitionistic linear logic) to Kreisel's modified realizability, the Diller-Nahm interpretation, and Gödel's Dialectica interpretation, as:

$ A _a^x$	Interpretation of IL^ω
$! \forall y A _y^x$	Kreisel modified realizability
$! \forall y \in a A _y^x$	Diller-Nahm interpretation
$! A _a^x$	Gödel's Dialectica interpretation.

But first we introduce a simplified version of the translation $(\cdot)^*$ from IL^ω into ILL_r^ω , which we will use in the treatment of the Diller-Nahm and the Dialectica interpretations (for modified realizability we use the translation $(\cdot)^\circ$). This simplification of Girard's translation is necessary so as to obtain an exact match between intuitionistic and linear interpretations. The simplification, however, requires two additional principles which, as we will see, turn out to be interpretable.

Proposition 6. *Consider the following simplification of Girard's translation $(\cdot)^*$, where the translation of \vee and \exists no longer needs the introduction of $!$ (cf. Proposition 1)*

$$\begin{aligned}
 A_{\text{at}}^+ &::= A_{\text{at}}, \quad \text{if } A_{\text{at}} \not\equiv \perp \\
 \perp^+ &::= 0 \\
 (A \wedge B)^+ &::= A^+ \& B^+ \\
 (A \vee B)^+ &::= A^+ \oplus B^+ \\
 (A \rightarrow B)^+ &::= !A^+ \multimap B^+ \\
 (\forall xA)^+ &::= \forall xA^+ \\
 (\exists xA)^+ &::= \exists xA^+.
 \end{aligned}$$

If A is provable in IL^ω then A^+ is provable in $\text{ILL}_r^\omega + \text{P}_\oplus + \text{P}_\exists$, where

$$\text{P}_\oplus : !(A \oplus B) \multimap !A \oplus !B$$

$$\text{P}_\exists : !\exists xA \multimap \exists x!A.$$

Proof. First we show that given the principles P_\oplus and P_\exists , we have $!A^* \multimap !A^+$. The proof is done by induction on the complexity of the formula A . Conjunction, implication and universal quantification follow easily by induction hypothesis using that ILL_r^ω proves:

$$\begin{aligned}
 !(A \& B) &\multimap !A \otimes !B \\
 !(A \multimap B) &\multimap !(A \multimap !B) \\
 !\forall xA &\multimap !\forall x!A
 \end{aligned}$$

respectively. Disjunction and existential quantification are studied below:

$$\begin{aligned}
 !(A \vee B)^* &\equiv !(A^* \oplus B^*) \multimap !A^* \oplus !B^* \\
 &\stackrel{(\text{IH})}{\multimap} !A^+ \oplus !B^+ \stackrel{(\text{P}_\oplus)}{\multimap} !(A^+ \oplus B^+) \equiv !(A \vee B)^+
 \end{aligned}$$

and $!(\exists xA)^* \equiv !\exists x!A^* \multimap \exists x!A^* \stackrel{(\text{IH})}{\multimap} \exists x!A^+ \stackrel{(\text{P}_\exists)}{\multimap} !\exists xA^+ \equiv !(\exists xA)^+$. Applying Proposition 1, we know that from $\text{IL}^\omega \vdash A$ we have $\text{ILL}_r^\omega \vdash A^*$. So, $\text{ILL}_r^\omega \vdash !A^*$ and hence $\text{ILL}_r^\omega + \text{P}_\oplus + \text{P}_\exists \vdash !A^*$. Using the equivalence proved before we have $\text{ILL}_r^\omega + \text{P}_\oplus + \text{P}_\exists \vdash !A^+$. In particular, we conclude $\text{ILL}_r^\omega + \text{P}_\oplus + \text{P}_\exists \vdash A^+$. \square

The reason we are freely allowed to assume the principles P_\oplus and P_\exists is that they are interpretable in all choices of interpretations we consider. Let us argue that P_\oplus and P_\exists are interpretable, by showing that the interpretation of premise implies that of the conclusion (hence the identity and projection functions can be taken as realisers for the implication). For the three choices of $\forall x \sqsubset a A$ we have considered one can show that

$$\begin{aligned}
 \forall x \sqsubset a (A(x) \& B) &\multimap (\forall x \sqsubset a A(x) \& B) \text{ and} \\
 \forall x \sqsubset a (B \multimap A(x)) &\multimap (B \multimap \forall x \sqsubset a A(x))
 \end{aligned}$$

when the variable x does not occur free in B . Also, $!(A \diamond_b B) \multimap !A \diamond_b !B$. Therefore, we have that

$$\begin{aligned}
|!(A \oplus B)|_{a,c}^{x,v,b} &\equiv !\forall y \sqsubset a \forall w \sqsubset c (|A|_y^x \diamond_b |B|_w^v) \\
&\rightarrow !(\forall y \sqsubset a |A|_y^x \diamond_b \forall w \sqsubset c |B|_w^v) \\
&\rightarrow !\forall y \sqsubset a |A|_y^x \diamond_b !\forall w \sqsubset c |B|_w^v \equiv |!A \oplus !B|_{a,c}^{x,v,b}.
\end{aligned}$$

Similarly, $|!\exists z A|_a^{x,z} \equiv !\forall y \sqsubset a |\exists z A|_y^{x,z} \equiv !\forall y \sqsubset a |A|_y^x \equiv |!A|_a^x \equiv |\exists z !A|_a^{x,z}$. Therefore, we can make use of the principles P_\oplus and P_\exists to simplify the embeddings of intuitionistic logic into (this extension of) linear logic, since the interpretation of linear logic will interpret these principles taking us back to standard linear logic (without P_\oplus and P_\exists). This is illustrated in the following diagram, where ILL_p^ω abbreviates $\text{ILL}_r^\omega + P_\oplus + P_\exists$ and $\text{LL}_{\text{ef}}^\omega$ abbreviates LL^ω without disjunctions and existential quantifications:

$$\begin{array}{ccc}
\text{ILL}_p^\omega & \xrightarrow{\quad} & \text{ILL}_b^\omega \\
\uparrow (\cdot)^+ & \text{|\cdot|} & \uparrow (\cdot)^+ = (\cdot)^* \\
\text{LL}^\omega & \xrightarrow{\text{Interpretation}} & \text{LL}_{\text{ef}}^\omega
\end{array}$$

The equality on the rightmost upward arrow represents the fact that all our interpretations transform proofs in LL^ω into proofs in $\text{LL}_{\text{ef}}^\omega$, where the two translations $(\cdot)^*$ and $(\cdot)^+$ coincide.

4.1 Modified realizability

Kreisel's modified realizability associates with each formula A of intuitionistic logic a new formula " $\mathbf{x} \text{ mr } A$ ", where \mathbf{x} is a sequence of fresh variables not present in A (see [18] for the formal definition). We are going to prove that this form of realizability once translated to the linear logic context via $(\cdot)^\circ$ corresponds (according to Theorem 2 below) to the interpretation of ILL_r^ω with $|!A|_y^x := !\forall y |A|_y^x$. First an auxiliary result:

Lemma 2. $|A^\circ|_y^x \multimap |!A|_y^x$.

Proof. Note that, because of the way we interpret $!A$, it can be checked by induction on A that the interpretation of A° has an empty tuple of challenge variables, i.e. we obtain a formula of the form $|A^\circ|_y^x$. To verify the lemma, it is enough to prove that $|A^\circ|_y^x \multimap !A'$, for some formula A' , since assuming this we have $|!A|_y^x \multimap !A' \multimap !A' \multimap |A^\circ|_y^x$. The proof is done by induction on the complexity of the formula A . We just sketch the cases of conjunction and disjunction, the other cases being immediate.

$$\begin{aligned}
 |(A \wedge B)^\circ|^{x,y} &\equiv |A^\circ \otimes B^\circ|^{x,y} \\
 &\equiv |A^\circ|^x \otimes |B^\circ|^y \\
 &\stackrel{(IH)}{\circ\text{-}\circ} !A' \otimes !B' \circ\text{-}\circ !(A' \& B'). \\
 |(A \vee B)^\circ|^{x,y,z} &\equiv |A^\circ \oplus B^\circ|^{x,y,z} \equiv |A^\circ|^x \diamond_z |B^\circ|^y \\
 &\stackrel{(IH)}{\circ\text{-}\circ} !A' \diamond_z !B' \\
 &\stackrel{(L1(v))}{\circ\text{-}\circ} !(A' \diamond_z B').
 \end{aligned}$$

That other cases are treated similarly. \square

Theorem 2. $|A^\circ|^x \circ\text{-}\circ (x \text{ mr } A)^\circ$.

Proof. The proof is done by induction on the complexity of the formula A . If A is an atomic formula, the result is trivial. Consider the case of conjunction:

$$\begin{aligned}
 |(A \wedge B)^\circ|^{x,y} &\equiv |A^\circ \otimes B^\circ|^{x,y} \equiv |A^\circ|^x \otimes |B^\circ|^y \\
 &\stackrel{(IH)}{\circ\text{-}\circ} (x \text{ mr } A)^\circ \otimes (y \text{ mr } B)^\circ \\
 &\equiv (x \text{ mr } A \wedge y \text{ mr } B)^\circ \equiv (x, y \text{ mr } A \wedge B)^\circ.
 \end{aligned}$$

The universal and existential quantifications also follow immediately using the induction hypothesis, and the way we define the translation and the interpretations. Implication is treated as

$$\begin{aligned}
 |(A \rightarrow B)^\circ|^g &\equiv |!(A^\circ \multimap B^\circ)|^g \equiv !\forall x |A^\circ \multimap B^\circ|^g_x \\
 &\equiv !\forall x (|A^\circ|^x \multimap |B^\circ|^g_x) \\
 &\stackrel{(IH)}{\circ\text{-}\circ} !\forall x ((x \text{ mr } A)^\circ \multimap (g x \text{ mr } B)^\circ) \\
 &\circ\text{-}\circ !\forall x !(x \text{ mr } A)^\circ \multimap (g x \text{ mr } B)^\circ \\
 &\equiv (\forall x (x \text{ mr } A \rightarrow g x \text{ mr } B))^\circ \equiv (g \text{ mr } (A \rightarrow B))^\circ
 \end{aligned}$$

whereas disjunction uses the auxiliary result above:

$$\begin{aligned}
 |(A \vee B)^\circ|^{x,y,z} &\stackrel{(L2)}{\circ\text{-}\circ} |!(A \vee B)^\circ|^{x,y,z} \equiv |!A^\circ \oplus B^\circ|^{x,y,z} \equiv |!(A^\circ|^x \diamond_z |B^\circ|^y)| \\
 &\stackrel{(IH)}{\circ\text{-}\circ} !((!(z = T) \multimap (x \text{ mr } A)^\circ) \& (!(z = F) \multimap (y \text{ mr } B)^\circ)) \\
 &\circ\text{-}\circ !((!(z = T) \multimap (x \text{ mr } A)^\circ) \otimes (!(z = F) \multimap (y \text{ mr } B)^\circ)) \\
 &\equiv ((z = T \rightarrow x \text{ mr } A) \wedge (z = F \rightarrow y \text{ mr } B))^\circ \\
 &\equiv (x, y, z \text{ mr } A \vee B)^\circ.
 \end{aligned}$$

That concludes the proof. \square

4.2 Gödel's Dialectica interpretation

Recall that Gödel's Dialectica interpretation first associates with each formula A a quantifier-free formula $A_D(\mathbf{x}; \mathbf{y})$ inductively. Then, each formula A is interpreted as the new formula $\exists \mathbf{x} \forall \mathbf{y} A_D(\mathbf{x}; \mathbf{y})$ (see [1], section 2.3). The next result shows that the Dialectica interpretation corresponds to the ILL_r^ω interpretation where $!A|_y^x := !A|_y^x$, via the simplified embedding $(\cdot)^+$ (cf. Proposition 6).

Theorem 3. $|A^+|_y^x \circ\text{-}\circ (A_D(\mathbf{x}; \mathbf{y}))^+$.

Proof. The proof is again an easy induction on the complexity of the formula A . The atomic formulas are checked trivially and the other formulas follow immediately by induction hypothesis using the definitions of the $(\cdot)^+$ -translation and the interpretations. We illustrate with two cases: conjunction

$$\begin{aligned} |(A \wedge B)^+|_{y,w}^{x,v} &\equiv |A^+ \& B^+|_{y,w}^{x,v} \equiv |A^+|_y^x \& |B^+|_w^v \\ &\stackrel{(\text{IH})}{\circ\text{-}\circ} (A_D(\mathbf{x}; \mathbf{y}))^+ \& (B_D(\mathbf{v}; \mathbf{w}))^+ \\ &\equiv (A_D(\mathbf{x}; \mathbf{y}) \wedge B_D(\mathbf{v}; \mathbf{w}))^+ \equiv ((A \wedge B)_D(\mathbf{x}, \mathbf{v}; \mathbf{y}, \mathbf{w}))^+ \end{aligned}$$

and disjunction

$$\begin{aligned} |(A \vee B)^+|_{y,w}^{x,v,z} &\equiv |A^+ \oplus B^+|_{y,w}^{x,v,z} \equiv |A^+|_y^x \diamond_z |B^+|_w^v \\ &\equiv (! (z = \text{T}) \text{-}\circ |A^+|_y^x) \& (! (z = \text{F}) \text{-}\circ |B^+|_w^v) \\ &\stackrel{(\text{IH})}{\circ\text{-}\circ} (! (z = \text{T}) \text{-}\circ (A_D(\mathbf{x}; \mathbf{y}))^+) \& (! (z = \text{F}) \text{-}\circ (B_D(\mathbf{v}; \mathbf{w}))^+) \\ &\equiv (z = \text{T} \rightarrow A_D(\mathbf{x}; \mathbf{y}))^+ \& (z = \text{F} \rightarrow B_D(\mathbf{v}; \mathbf{w}))^+ \\ &\equiv ((z = \text{T} \rightarrow A_D(\mathbf{x}; \mathbf{y})) \wedge (z = \text{F} \rightarrow B_D(\mathbf{v}; \mathbf{w})))^+ \\ &\equiv ((A \vee B)_D(\mathbf{x}, \mathbf{v}, z; \mathbf{y}, \mathbf{w}))^+. \end{aligned}$$

The other cases are treated similarly. \square

Note that although $(\cdot)^+$ translates formulas from IL^ω into $\text{ILL}_r^\omega + \text{P}_\oplus + \text{P}_\exists$, since these two principles are interpretable the verifying system is still ILL_b^ω .

4.3 Diller-Nahm interpretation

The Diller-Nahm interpretation differs from Gödel's Dialectica interpretation since it allows finite sets to witness the negative content of an implication. Formally, the Diller-Nahm interpretation can be defined inductively as

$$\begin{aligned} (A_{\text{at}})_{dn}(\cdot) &:\equiv A_{\text{at}} \\ (A \wedge B)_{dn}(\mathbf{x}, \mathbf{v}; \mathbf{y}, \mathbf{w}) &:\equiv A_{dn}(\mathbf{x}; \mathbf{y}) \wedge B_{dn}(\mathbf{v}; \mathbf{w}) \\ (A \vee B)_{dn}(\mathbf{x}, \mathbf{v}, z; \mathbf{y}, \mathbf{w}) &:\equiv (z = \text{T} \rightarrow A_{dn}(\mathbf{x}; \mathbf{y})) \wedge (z = \text{F} \rightarrow B_{dn}(\mathbf{v}; \mathbf{w})) \\ (A \rightarrow B)_{dn}(f, g; \mathbf{x}, \mathbf{w}) &:\equiv \forall y \in f \mathbf{xw} A_{dn}(\mathbf{x}; \mathbf{y}) \rightarrow B_{dn}(g\mathbf{x}; \mathbf{w}) \\ (\forall z A)_{dn}(f; \mathbf{y}, z) &:\equiv A_{dn}(fz; \mathbf{y}) \\ (\exists z A)_{dn}(\mathbf{x}, z; \mathbf{y}) &:\equiv A_{dn}(\mathbf{x}; \mathbf{y}). \end{aligned}$$

Next we show that the Diller-Nahm interpretation of IL^ω corresponds to the interpretation of ILL_r^ω with $!A|_a^x := !\forall y \in a |A|_y^x$.

Theorem 4. $|A^+|_y^x \multimap (A_{dn}(x; y))^+$.

Proof. The proof, by induction on the structure of A , is similar to the one concerning Gödel's interpretation. The only case which needs attention is that of implication, which we analyse below.

$$\begin{aligned}
 |(A \rightarrow B)^+|_{x,w}^{f,g} &\equiv !A^+ \multimap B^+|_{x,w}^{f,g} \equiv !A^+|_{fxw}^x \multimap |B^+|_w^{gx} \\
 &\equiv !\forall y \in fxw |A^+|_y^x \multimap |B^+|_w^{gx} \\
 &\stackrel{\text{(IH)}}{\multimap} !\forall y \in fxw (A_{dn}(x; y))^+ \multimap (B_{dn}(gx; w))^+ \\
 &\equiv !(\forall y \in fxw A_{dn}(x; y))^+ \multimap (B_{dn}(gx; w))^+ \\
 &\equiv (\forall y \in fxw A_{dn}(x; y) \rightarrow B_{dn}(gx; w))^+ \\
 &\equiv ((A \rightarrow B)_{dn}(f, g; x, w))^+.
 \end{aligned}$$

Note that the $(\cdot)^+$ translation of $\forall y \in a A$ is $\forall y \in a A^+$, as we can see below:

$$\begin{aligned}
 (\forall y \in a A)^+ &\equiv (\forall y (y \in a \rightarrow A))^+ \\
 &\equiv \forall y (!y \in a \multimap A^+) \equiv \forall y (!y \in a) \multimap A^+ \equiv \forall y \in a A^+.
 \end{aligned}$$

That concludes the proof. \square

5 The Categorical Approach

The study developed in this paper (and in previous work of the second author) is strongly inspired by work of de Paiva and Hyland on categorical models of linear logic using Gödel's Dialectica interpretation. In this section we try to explain and make more explicit the link between our framework for unifying interpretations of IL via interpretations of ILL and the categorical approach on [14–16] for modelling ILL . More precisely, in [14] one finds a categorical version of the Dialectica interpretation and an endofunctor interpretation for the modality $!A$ that corresponds to the Diller-Nahm interpretation. Our goal is to relate this approach with the work in the previous sections.

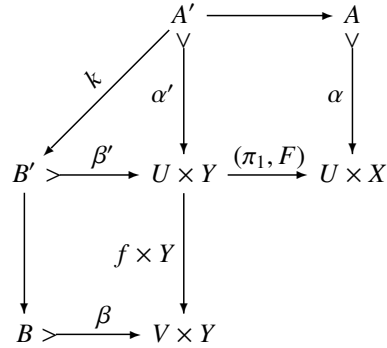
Before presenting de Paiva's category DC that models ILL , for sake of intuition, let us informally sketch the correspondence between our framework and hers through the following table.

	Our framework	de Paiva's framework
Realizers in	\mathcal{T}^ω - finite types	\mathbf{C} - cartesian closed category
Formulas	$ A \subseteq X \times Y$	$X \stackrel{\alpha}{\leftarrow} Y$ (object of DC)
Sequents	$A \vdash B$	$A \xrightarrow{(f,F)} B$ (morphism of DC)
Linear implication	$A \multimap B$	$[A, B]_{\text{DC}}$ or B^A

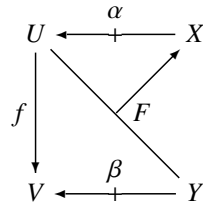
First, we point out that in de Paiva’s work the realisers of the functional interpretation are taken from a given (fixed) cartesian closed category \mathbf{C} . In our case, we work with the particular cartesian closed category of the functionals of finite type. Also, our interpretations are given syntactically, and hence, a formula A is interpreted as another formula $|A|_y^x$, which can be thought of as a binary relation between x and y . In de Paiva’s work these relations are at the core of constructing a new category \mathbf{DC} out of the given ccc \mathbf{C} .

Let us briefly describe how the category \mathbf{DC} is defined and its associated constructions. Starting with \mathbf{C} , a finitely complete cartesian closed category with stable and disjoint coproducts, we can define the monoidal closed category \mathbf{DC} as follows. An object of \mathbf{DC} is a subobject of the product $U \times X$, thus a monomorphism $A \xrightarrow{\alpha} U \times X$ with A, U and X objects of \mathbf{C} also denoted by $U \stackrel{\alpha}{\leftarrow} X$. If we think of these objects as set-theoretic relations between U and X , and considering α as the identity monic, we get that $A \subseteq U \times X$, precisely as in our framework.

A map between two such objects $A \xrightarrow{\alpha} U \times X$ and $B \xrightarrow{\beta} V \times Y$ consists of a pair of maps of \mathbf{C} , (f, F) , $f : U \rightarrow V$, $F : U \times Y \rightarrow X$ such that pulling back $A \xrightarrow{\alpha} U \times X$ along $U \times Y \xrightarrow{(\pi_1, F)} U \times X$ and $B \xrightarrow{\beta} V \times Y$ along $U \times Y \xrightarrow{f \times Y} V \times Y$ (see the diagram below), the first subobject $A' \xrightarrow{\alpha'} U \times Y$ is smaller than the second $B' \xrightarrow{\beta'} U \times Y$, i.e. there is a map $k : A' \rightarrow B'$ in \mathbf{C} making the triangle in the diagram below commute:



If we write the two relations in the short version $U \stackrel{\alpha}{\leftarrow} X$ and $V \stackrel{\beta}{\leftarrow} Y$ and $(-)^{-1}$ for the pullback functor, then a map in \mathbf{DC} can be represented as the pair (f, F) in the diagram below



satisfying the condition $(\pi_1, F)^{-1}(\alpha) \leq (f \times Y)^{-1}(\beta)$.

The intuition in terms of set-theoretic relations is the following: there is a map $\alpha \xrightarrow{(f,F)} \beta$ in DC if and only if whenever $u\alpha F(u, y)$ then $f(u)\beta y$. In what follows we are going to say that two elements, x and y are related via α (i.e. $x\alpha y$) by α_x^x . This way the comparison with our framework becomes easier. Using this notation, the condition above says that whenever $\alpha_{F(u,y)}^u$ then $\beta_y^{f(u)}$.

In the category DC we can also define the bifunctors \otimes , $[-, -]_{\text{DC}}$ and $\&$ and the operation \oplus of weak-coproducts that can be read intuitively as

$$\begin{aligned} (\alpha \otimes \beta)_{x,y}^{u,v} & \text{ iff } \alpha_x^u \text{ and } \beta_y^v \\ [\alpha, \beta]_{v,z}^{F,f} \equiv (\beta^\alpha)_{v,z}^{F,f} & \text{ iff } \alpha_{F(v,z)}^v \Rightarrow \beta_z^{f(v)} \\ (\alpha \& \beta)_w^{u,v} & \text{ iff } \alpha_w^u \text{ or } \beta_w^v \text{ depending whether } w \text{ is in } X \text{ or } Y \\ (\alpha \oplus \beta)_{f,g}^w & \text{ iff } \alpha_{f(w)}^w \text{ or } \beta_{g(w)}^w \text{ depending whether } w \text{ is in } U \text{ or } V. \end{aligned}$$

Apart from the relation $\alpha \otimes \beta$, our interpretation of the linear logic connectives (Definition 1) coincides precisely with the definitions above. Let us examine in more detail the interpretation of \oplus , where our two approaches lead to different interpretations.

The main reason why we can have a simpler definition of $\alpha \oplus \beta$ (with no need for the second player to play higher order moves f, g) is because we always assume that each finite type is inhabited by at least one element, while de Paiva's imposes no similar restriction. More precisely, in our setting we have

$$|\alpha \oplus \beta| \subseteq (U \times V \times \mathbb{B}) \times (X \times Y),$$

with \mathbb{B} for the set of boolean constants, while in de Paiva's setting, considering set-theoretic relations,

$$(\alpha \oplus \beta) \subseteq (U + V) \times (X^U \times Y^V).$$

If U and V are non-empty, then the two types $U \times V \times \mathbb{B}$ and $U + V$ are isomorphic. In case, however, one of U or V is empty then $U \times V \times \mathbb{B}$ is also empty, whereas $U + V$ can still be non-empty. In other words, in the most general case, when types can be empty, we must indeed work with the type $U + V$ rather than with $U \times V \times \mathbb{B}$. Let us see then, how the interpretation of $\alpha \oplus \beta$ works in the case when some of the move-sets of Eloise could be empty.

While in the first situation Eloise plays one element of U , one from V , and a boolean choosing which game is going to count, in the second case Eloise plays an element of $U + V$. As we are going to see, in the latter case (with no extra assumptions) we need Abelard to play functions. Consider the \oplus L-rule, where from $\Gamma, A \vdash C$ and $\Gamma, B \vdash C$ we can conclude $\Gamma, A \oplus B \vdash C$ (for simplicity we shall omit the context Γ). In our framework, the proof of the two premises will provide realisers F and H such that the premises of the following rule are derivable:

$$\frac{|A|_{Fu}^u \multimap |C|_w^{fu} \quad |B|_{Hv}^v \multimap |C|_w^{hv}}{|A \oplus B|_{Fu, Hv}^{u,v,b} \multimap |C|_w^{b(fu, hv)}}.$$

In order to realise C in the conclusion of the rule we can make use of the boolean b to choose between fu and hv . Moreover, the “negative” realiser for $A \oplus B$ is just a pair

$\langle Fuw, Hvw \rangle$. Now, in de Paiva’s more general setting, we have the same information about the realisers for the premises of the rule

$$\frac{|A|_{Fuw}^u \multimap |C|_w^{fu} \quad |B|_{Hvw}^v \multimap |C|_w^{hv}}{|A \oplus B|_{F'aw, H'aw}^a \multimap |C|_w^{(a \in U)(fa)(ha)}}.$$

but in the conclusion we no longer have a triple $\langle u, v, b \rangle$, but rather an element $a \in U + V$. Therefore, the functionals F and H are lifted to functionals F' and H' as

$$F'aw := \begin{cases} \lambda u^U. Faw & \text{if } a \in U \\ \lambda u^U. Faw & \text{if } a \in V \end{cases} \quad H'aw := \begin{cases} \lambda v^V. Hvw & \text{if } a \in U \\ \lambda v^V. Hvw & \text{if } a \in V. \end{cases}$$

The extra arguments u and v are used in the cases when the parameter a has the “wrong” type to be used in either F or G , and a standard value must be used.

For the rest of the section, let us analyse how the linear logic exponential $!A$ is interpreted in both approaches. As pointed by Blute and Scott in [3], apropos natural and satisfying categorical models for the LL connectives “*unfortunately, the exponentials are less clear: the structure seems less canonical*”. In terms of monoidal categories the structure used to model $!A$ is that of comonads and comonoid objects. In [14], it is shown that if the category \mathbf{C} has a free monoid structure with countable coproducts then the endofunctor $!$ can be defined on objects of \mathbf{DC} as the pullback of $A^* \xrightarrow{\alpha^*} (U \times X)^*$ along $U \times X^* \xrightarrow{C(U,X)} (U \times X)^*$:

$$\begin{array}{ccccc} !A & \longrightarrow & A^* & & A \\ \downarrow \text{!}\alpha & & \downarrow \alpha^* & \longleftarrow * & \downarrow \alpha \\ U \times X^* & \xrightarrow{C(U,X)} & (U \times X)^* & & U \times X \end{array}$$

Note that the functor $*$: $\mathbf{C} \rightarrow \mathbf{MonC}$ is left-adjoint to the forgetful functor $U : \mathbf{MonC} \rightarrow \mathbf{C}$ (see [14] for more details). Intuitively, the relation α_x^u is transformed into a new relation $(!\alpha)_{\{x_1, \dots, x_n\}}^u$ which is equivalent to $\forall x \in \{x_1, \dots, x_n\} \alpha_x^u$. The functor $!$ acts on morphisms in \mathbf{DC} as $!(f, F) \equiv (f, !F)$ where $!F : U \times Y^* \rightarrow X^*$ is the composite of

$$U \times Y^* \xrightarrow{C(U,Y)} (U \times Y)^* \xrightarrow{F^*} X^*.$$

Since the functor $! : \mathbf{DC} \rightarrow \mathbf{DC}$ has a natural comonad $(!, \epsilon, \delta)$ structure and $!A$ is a comonoid object in \mathbf{DC} , $!$ models the linear logic exponential in the style of the Diller-Nahm variant of the Dialectica interpretation, via finite sets.

In our approach, we have chosen to take a formal (syntactic) approach for the interpretation of $!A$. We identify three conditions (A1–A3) which $!A$ needs to satisfy in order for the resulting interpretation to be sound. Our conditions are more general, and include as particular case the instance where $!$ is a comonad with comonoid objects. In particular, we are able to obtain interpretations of $!A$ that correspond to other well-known functional interpretation such as Gödel’s Dialectica interpretation and Kreisel’s

modified realizability. A natural question, of course, arises: Do the Dialectica and modified realizability interpretations fit into the framework of de Paiva as well, and can they be seen as arising from other comonads with comonoidal structure? In the first case (the Dialectica interpretation) the answer is yes, and de Paiva does have a few remarks about the Dialectica interpretation in her paper and in her thesis. More precisely, let $! : \text{DC} \rightarrow \text{DC}$ be the identity endofunctor. Intuitively $(!\alpha)_v^u$ if and only if α_v^u . It is immediate to check that $(!, id, id)$ is a comonad, but in order for $!\alpha$ to be a comonoid object in DC (not surprisingly) we need to require decidability. More precisely, $!\alpha \rightarrow (!\alpha \otimes !\alpha)$ is interpreted as $(!\alpha)_{y_0, y_1}^x \rightarrow (!\alpha)_{y_0}^x \otimes (!\alpha)_{y_1}^x$ with

$$y_0 \cdot y_1 := \begin{cases} y_0 & \text{if } \neg \alpha_{y_0}^x \\ y_1 & \text{otherwise.} \end{cases}$$

As for modified realizability, it is not clear to us at the moment whether it can also be shown to arise from a different monoid (other than the free monoid) using a generalisation of de Paiva's construction. We plan to consider this question in our future investigations.

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