

Recent Developments in Proof Mining

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Outline

- 1 Introduction
 - Proof Mining
 - Functional Interpretations

- 2 Recent Case Studies
 - Approximation Theory
 - Fixed Point Theory
 - Ergodic Theory

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*Extraction of computational content from
(ineffective) mathematical proofs*



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*Proofs often carry more information than
what is stated as theorem*



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Main Technique

Functional interpretations:

- **Dialectica** (Gödel'1958)
- **Diller-Nahm variant** (Diller/Nahm'1974)
- **Monotone Dialectica** (Kohlenbach'1990)
- **Bounded Dialectica** (Ferreira/O.'2005)



Simple Example

Theorem

$$\forall n(f(n+1) > f(n)) \rightarrow \forall k(f(k+2) > f(k))$$



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Can we compute n given k ?



Dialectica Interpretation

Interpret

$$\forall n(f(n+1) > f(n)) \rightarrow \forall k(f(k+2) > f(k))$$

as

$$\exists \phi \forall k(f(\phi k + 1) > f(\phi k) \rightarrow f(k+2) > f(k))$$



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Witness can be produced, e.g.

$$\phi k := \begin{cases} k & \text{if } f(k+1) \leq f(k) \\ k+1 & \text{otherwise} \end{cases}$$



Diller-Nahm Variant

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$$\forall n(f(n+1) > f(n)) \rightarrow \forall k(f(k+2) > f(k))$$

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Witness can be produced, e.g.

$$\phi k := \{k, k+1\}$$



Bounded Dialetica Interpretation

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$$\forall n(f(n+1) > f(n)) \rightarrow \forall k(f(k+2) > f(k))$$

as

$$\exists \phi \forall k (\forall m \leq \phi k (f(m+1) > f(m)) \rightarrow f(k+2) > f(k))$$



Bounded Diagonal Interpretation

Interpret

$$\forall n(f(n+1) > f(n)) \rightarrow \forall k(f(k+2) > f(k))$$

as

$$\exists \phi \forall k (\forall m \leq \phi k (f(m+1) > f(m))) \rightarrow f(k+2) > f(k)$$

Witness can be produced, e.g.

$$\phi k := k + 1$$



Monotone Dialectica Interpretation

Interpret

$$\forall n(f(n+1) > f(n)) \rightarrow \forall n(f(n+2) > f(n))$$

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$$\exists \phi \exists \psi \leq^* \phi \forall k(f(\psi k + 1) > f(\psi k) \rightarrow f(k+2) > f(k))$$



Monotone Dialectica Interpretation

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Bounds don't depend on bounded input

E.g. compact spaces



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- **Ineffective principles become interpretable**

Witnesses may not be computable but can be bounded

E.g. WKL



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- **Not much is lost**

Bounds often give precise witness

E.g. monotonicity, searchable set

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- Existence and uniqueness of best approximations
E.g. approximate continuous functions by polynomials



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- **Existence:** quite often ineffective, non-computational
- **Uniqueness:** of form that proof mining applies

$$\forall n^{\mathbb{N}}, f^{C[0,1]}, p_1^{P_n}, p_2^{P_n} (\|f - p_i\| = \text{best} \rightarrow \|p_1 - p_2\| = 0)$$



Approximation Theory

- Existence and uniqueness of best approximations
E.g. approximate continuous functions by polynomials
- **Existence:** quite often ineffective, non-computational
- **Uniqueness:** of form that proof mining applies

$$\forall n^{\mathbb{N}}, f^{C[0,1]}, p_1^{P_n}, p_2^{P_n} (\|f - p_i\| = \text{best} \rightarrow \|p_1 - p_2\| = 0)$$

$$\forall n, f, p_1, p_2, l \exists k (\|f - p_i\| - \text{best} \leq 2^{-k} \rightarrow \|p_1 - p_2\| < 2^{-l})$$



L_1 Approximation

Theorem (Jackson'1921)

For any fixed $n \in \mathbb{N}$ and continuous function $f \in C[0, 1]$ there exists a unique polynomial $p_n \in P_n$ such that $\|f - p_n\|_1$ is minimal.



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Proof (Cheney'1965).

Mathematically elementary proof (just 2 pages), but logically intricate.
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How to compute p_n given f and n ?

- Partial results during the 1970's
[Björnestål'1975 and Kroó'1978]
- Explicit algorithm extracted from Cheney's 1965 proof
[Kohlenbach/O. 2001]



Main Obstacle

Attainment of the infimum (WKL) used in proof of following lemma

Lemma (Original)

$\forall x \in A (f(x) \neq 0) \rightarrow \dots$



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WKL used to obtain distance from zero

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WKL used to obtain distance from zero

$$\forall x \in A (f(x) \neq 0) \rightarrow \exists \delta \forall x \in A (|f(x)| \geq \delta)$$

We showed that the weaker version of the lemma is sufficient

Lemma (Weakening)

$$\exists \delta \forall x \in A (|f(x)| \geq \delta) \rightarrow \dots$$



History (L_1 Approximation)

1921	Jackson	proof of existence and uniqueness
1965	Cheney	elementary proof of uniqueness
1975	Björnestrål	ineff. existence of modulus on f, n
1978	Kroó	ineff. existence of modulus on ω_f, n
2001	Kohlenbach/O.	explicit modulus of uniqueness
2002	Oliva	complexity of L_1 approximation



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Banach Theorem (1922)

- (X, d) complete metric space
- $f : X \rightarrow X$ is **contractive** if $d(f(x), f(y)) \leq \delta \cdot d(x, y)$ ($\delta < 1$)



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Theorem (Banach'1922)

If $f : X \rightarrow X$ is contractive then f has a unique fixed-point.



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Theorem (Banach'1922)

If $f : X \rightarrow X$ is contractive then f has a unique fixed-point.

For any $x_0 \in X$, $x_{n+1} := f(x_n)$ converges to the fixed-point



Browder/Göhde/Kirk Theorem (1965)

- $(X, \|\cdot\|)$ uniformly convex Banach space
- $C \subseteq X$ convex, closed and bounded
- $f : C \rightarrow C$ is **nonexpansive** if $\|f(x) - f(y)\| \leq \|x - y\|$



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Theorem (Browder, Göhde, Kirk'1965)

If $f : C \rightarrow C$ is nonexpansive then f has a fixed-point.

- If $f(C)$ compact $x_{n+1} := \frac{x_n + f(x_n)}{2}$ converges to a fixed-point
- Rate of convergence in general not computable (Kohlenbach)
- Can compute rate of **asymptotic regularity** of x_n
i.e. how fast $\|x_n - f(x_n)\| \rightarrow 0$



Ishikawa Theorem (1976)

- $(X, \|\cdot\|)$ normed linear space
- $C \subseteq X$ convex and bounded

Theorem (Ishikawa'1976)

If $f : C \rightarrow C$ is nonexpansive then $\lim_{n \rightarrow \infty} \|x_n - f(x_n)\| = 0$, where

- $(\lambda_n)_{n \in \mathbb{N}} \in [0, 1]$ is divergent in sum and $\limsup \lambda_n < 1$
- $x_{n+1} = (1 - \lambda_n)x_n + \lambda_n f(x_n)$



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Theorem (Borwein/Reich/Shafir'1992)

Dropping the boundedness assumption on C we still have

$$\lim_{n \rightarrow \infty} \|x_n - f(x_n)\| = r_C(f).$$



History (asymptotic regularity)

1976	Ishikawa	no uniformity
1978	Edels./O'Brien	ineff. uniformity in x_0 (fixed λ)
1983	Goebel/Kirk	ineff. uniformity in x_0 and f (*)
1990	Goebel/Kirk	conjecture no uniformity in C
1992	Bor./Rei./Sha.	generalisation of Ishikawa (*)
1996	Baillon/Bruck	full uniformity for fixed λ
2001	Kohlenbach	full uniformity
<hr/>		
2000	Kirk	uniformity on x_0, f (f direc. nonexp., fixed λ)
2003	Kohlen./Leust.	full uniformity, hyper. spc. and f direc. nonexp.



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Ergodic

Systems or processes with the property that, given sufficient time, they include or affect all points in a given space



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Systems or processes with the property that, given sufficient time, they include or affect all points in a given space

Such systems or processes can be represented statistically by a reasonably large selection of points



Ergodic Theory

- Let
 (X, Σ, μ) probability space
 $T : X \rightarrow X$ measure preserving transformation



Ergodic Theory

- Let

(X, Σ, μ) probability space

$T : X \rightarrow X$ measure preserving transformation

- T is **ergodic** if

$$\mu(A) \neq 0 \wedge \mu(X \setminus A) \neq 0 \implies \mu(A \Delta T^{-1}(A)) \neq 0$$



Ergodic Theory

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 (X, Σ, μ) probability space
 $T : X \rightarrow X$ measure preserving transformation
- T is **ergodic** if
 $\mu(A) \neq 0 \wedge \mu(X \setminus A) \neq 0 \implies \mu(A \Delta T^{-1}(A)) \neq 0$
- Study of **ergodic transformations**



Mean Ergodic Theorem (functional analysis)

- \mathcal{H} Hilbert space
- $T : \mathcal{H} \rightarrow \mathcal{H}$ nonexpansive linear operator (i.e. $\|Tf\| \leq \|f\|$)
- $S_n f := f + Tf + \dots + T^{n-1}f$
- $A_n f := \frac{S_n f}{n}$



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Theorem (von Neumann)

The sequence $A_n f$ converges.

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Theorem (von Neumann)

The sequence $A_n f$ converges.

What about rate of convergence?



Rate of Convergence

Convergence:

$$\forall \varepsilon^{\mathbb{Q}^*} \exists n^{\mathbb{N}} \forall m \geq n (\|A_m f - A_n f\| < \varepsilon)$$



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Not computable in general!



Rate of Convergence (n.c.i)

Look at the no-counterexample interpretation of

$$\forall \varepsilon^{\mathbb{Q}^+} \exists n^{\mathbb{N}} \forall m \geq n (\|A_m f - A_n f\| < \varepsilon)$$

i.e.

$$\forall \varepsilon^{\mathbb{Q}^+}, M \exists n^{\mathbb{N}} (M(n) \geq n \rightarrow \|A_m f - A_n f\| < \varepsilon)$$

or, equivalently

$$\forall \varepsilon^{\mathbb{Q}^+}, K \exists n^{\mathbb{N}} \forall m \in [n, K(n)] (\|A_m f - A_n f\| < \varepsilon)$$



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or, equivalently

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Classically equivalent! Computationally better!



Avigad/Gerhardy

$$\forall \varepsilon \in \mathbb{Q}_+^*, K \exists n \in \mathbb{N} \forall m \in [n, K(n)] (\|A_m f - A_n f\| < \varepsilon)$$

Extraction of bound n given ε and K

Uses elimination of monotone Skolem functions (due to Kohlenbach)

Summary

- Approximation theory
 - *modulus of uniqueness*
 - classical logic, weak König's Lemma
- Fixed point theory
 - *modulus of asymptotic regularity*
 - convergence of bounded monotone sequences
- Ergodic theory
 - *modulus of convergence* (n.c.i.)
 - convergence of bounded monotone sequences

